Compressive Sensing in Noise and the Role of Adaptivity

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Compressive Sensing in Noise



When (and how well) can we estimate x from the measurements y?

Nonadaptive Compressive Sensing

Stable Signal Recovery

Given
$$y = \Phi x + e$$
, find x

Typical (worst-case) guarantee: If Φ satisfies the RIP

$$\|\widehat{x} - x\|_2^2 \le C \|e\|_2^2$$

Even if $\Lambda = \operatorname{supp}(x)$ is provided by an oracle, the error can still be as large as $\|\widehat{x} - x\|_2^2 = \|e\|_2^2/(1-\delta)$.

Stable Signal Recovery: Part II

Suppose now that Φ satisfies

$$A(1-\delta)\|x\|_{2}^{2} \le \|\Phi x\|_{2}^{2} \le A(1+\delta)\|x\|_{2}^{2} \qquad \|x\|_{0} \le 2S$$

In this case our guarantee becomes

$$\left\| \widehat{x} - x \|_{2}^{2} \le \frac{C}{A} \|e\|_{2}^{2} \right\|$$

Unit-norm rows $\|\widehat{x} - x\|_2^2 \leq C \frac{N}{M} \|e\|_2^2$

Expected Performance

- Worst-case bounds can be pessimistic
- What about the *average* error?
 - assume e is white noise with variance σ^2

 $\mathbb{E}\left(\|e\|_2^2\right) = M\sigma^2$

- for oracle-assisted estimator

$$\mathbb{E}\left(\|\widehat{x} - x\|_2^2\right) \le \frac{S\sigma^2}{A(1-\delta)}$$

- if e is Gaussian, then for ℓ_1 -minimization

$$\mathbb{E}\left(\|\widehat{x} - x\|_{2}^{2}\right) \leq \frac{C'}{A}S\sigma^{2}\log N$$

White Signal Noise

What if our signal x is contaminated with noise?

$$y = \Phi(x+n) = \Phi x + \Phi n$$

Suppose Φ has orthogonal rows with norm equal to \sqrt{B} . If n is white noise with variance σ^2 , then Φn is white noise with variance $B\sigma^2$.

$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \leq C' \frac{B}{A} S \sigma^{2} \log N$$

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$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \leq C' \frac{N}{M} S \sigma^{2} \log N$$

 $SNR = 10 \log_{10} \left(\frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right) \longrightarrow \begin{array}{c} \text{3dB loss per octave} \\ \text{of subsampling} \end{array}$

Noise Folding



[Davenport, Laska, Treichler, Baraniuk - 2011]

Room For Improvement?

There exists matrices Φ (with unit-norm rows) such that for *any* (sparse) x we have

$$\mathbb{E} \|\widehat{x} - x\|_2^2 \le C \frac{N}{M} S \sigma^2 \log N.$$
$$y_i = \langle \phi_i, x \rangle + e_i$$
$$\bigstar$$
$$\phi_i \text{ and } x \text{ are almost orthogonal}$$

- We are using most of our "sensing power" to sense entries that aren't even there!
- Tremendous loss in signal-to-noise ratio (SNR)
- It's hard to imagine any way to avoid this...

Can We Do Better?

Via a better choice of Φ ? Via a better recovery algorithm?

If
$$y = \Phi x + e$$
 with $e \sim \mathcal{N}(0, \sigma^2 I)$, then there exists an x such that for any \hat{x} and any Φ
$$\mathbb{E}\left[\|\widehat{x}(\Phi x + e) - x\|_2^2\right] \geq C \frac{N}{\|\Phi\|_F^2} S \sigma^2 \log(N/S).$$

If
$$y = \Phi(x + n)$$
 with $n \sim \mathcal{N}(0, \sigma^2 I)$, then there exists an x such that for any \hat{x} and any Φ

$$\mathbb{E}\left[\|\widehat{x}(\Phi(x+n)) - x\|_2^2\right] \ge C\frac{N}{M}S\sigma^2\log(N/S).$$

$$\Phi = U\Sigma V^* \quad y' = \Sigma^{-1} U^* y = V^* x + V^* n \quad \|V^*\|_F^2 = M$$
[Candès and Davenport - 2011]

Intuition

Suppose that y = x + n with $n \sim \mathcal{N}(0, I)$ and that S = 1

$$\mathbb{E} \|\widehat{x}(y) - x\|_2^2 \ge C' \log N$$



Proof Recipe

Ingredients (Makes $\sigma^2 = 1$ servings)

• Lemma 1: There exists a set \mathcal{X} of S-sparse vectors such that

•
$$|\mathcal{X}| = (N/S)^{S/4}$$

• $||x_i - x_j||_2 \ge \frac{1}{2}$ for all $x_i, x_j \in \mathcal{X}$
• $||\frac{1}{|\mathcal{X}|} \sum_i x_i x_i^* - \frac{1}{N}I|| \le \frac{\beta}{N}$ for some $\beta > 0$

• Lemma 2: Define $R^*_{mm}(\Phi) = \inf_{\widehat{x}} \sup_{\|x\|_0 \le S} \mathbb{E} \left[\|\widehat{x}(\Phi x + e) - x\|_2^2 \right].$

Suppose \mathcal{X} is a set of S-sparse vectors such that $\|x_i - x_j\|_2^2 \ge 8NR_{\min}^*(\Phi)$ for all $x_i, x_j \in \mathcal{X}$. Then $\frac{1}{2} \log |\mathcal{X}| - 1 \le \frac{1}{2|\mathcal{X}|^2} \sum_{i,j} \|\Phi x_i - \Phi x_j\|_2^2$.

Instructions

Combine ingredients and add a dash of linear algebra.

The Details

$$\mu = \frac{1}{|\mathcal{X}|} \sum_{i} x_{i} \quad Q = \frac{1}{|\mathcal{X}|} \sum_{i} x_{i} x_{i}^{*}$$

 $\frac{S}{4} \log(N/S) - 2 \le \frac{1}{|\mathcal{X}|^2} \sum_{i,j} \|\Phi x_i - \Phi x_j\|_2^2$ $= \operatorname{Tr}\left(\Phi^*\Phi\left(\frac{1}{|\mathcal{X}|^2}\sum_{i,j}(x_i - x_j)(x_i - x_j)^*\right)\right)$ $= \operatorname{Tr} \left(\Phi^* \Phi \left(2(Q - \mu \mu^*) \right) \right)$ $\leq 2 \operatorname{Tr} \left(\Phi^* \Phi Q \right)$ $\leq 2 \operatorname{Tr} (\Phi^* \Phi) \| Q \|$ $< 2 \|\Phi\|_{F}^{2} \cdot 16 R_{mm}^{*}(\Phi)(1+\beta)$ $R_{\rm mm}^*(\Phi) \ge \frac{S \log(N/S)}{128(1+\beta) \|\Phi\|^2}$

Lemma 1

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Strategy

Construct ${\mathcal X}$ by sampling (with replacement) from

$$\mathcal{U} = \left\{ x \in \{0, \sqrt{1/S}, -\sqrt{1/S}\}^n : \|x\|_0 \le S \right\}$$

Repeat for $|\mathcal{X}| = (N/S)^{S/4}$ iterations.

With probability > 0, the remaining properties are satisfied.

Key: Matrix Bernstein Inequality [Ahlswede and Winter, 2002]

Adaptive Sensing

Adaptive Sensing

Think of sensing as a game of 20 questions



Simple strategy: Use M/2 measurements to find the support, and the remainder to estimate the values.

Thought Experiment

Suppose that after M/2 measurements we have perfectly estimated the support.



Does Adaptivity *Really* Help?

Sometimes...

- Noise-free measurements, but non-sparse signal
 - adaptivity doesn't help if you want a uniform guarantee
 - probabilistic adaptive algorithms can reduce the required number of measurements from $O(S\log(N/S))$ to $O(S\log\log(N/S))$ [Indyk et al. 2011]
- Noisy setting
 - distilled sensing [Haupt et al. 2007, 2010]
 - adaptivity can reduce the estimation error to

Which is it?

Which Is It?

Suppose we have a budget of M measurements of the form $y_i = \langle \phi_i, x \rangle + e_i$ where $\|\phi_i\|_2 = 1$ and $e_i \sim \mathcal{N}(0, \sigma^2)$

The vector ϕ_i can have an arbitrary dependence on the measurement history, i.e., $(\phi_1, y_1), \ldots, (\phi_{i-1}, y_{i-1})$

Theorem

There exist x with $||x||_0 \le S$ such that for *any* adaptive measurement strategy and *any* recovery procedure \hat{x} ,

$$\mathbb{E} \|\widehat{x}(y) - x\|_2^2 \ge C \frac{N}{M} S \sigma^2.$$

Thus, in general, adaptivity does not seem to help!

[Arias-Castro, Candès, and Davenport - 2011]

Proof Strategy

- Step 1: Consider a prior on sparse signals with nonzeros of amplitude $\mu\approx\sigma\sqrt{N/M}$
- Step 2: Show that if given a budget of M measurements, you cannot detect the support very well
- Step 3: Immediately translate this into a lower bound on the MSE

To make things simpler, we will consider a Bernoulli prior $\pi(x)$ instead of a uniform S-sparse prior:

$$x_j = \begin{cases} 0 & \text{with probability } 1 - S/N \\ \mu > 0 & \text{with probability } S/N \end{cases}$$

Proof of Main Result

Let $T = \{j : x_j \neq 0\}$ and set $\sigma^2 = 1$ For any estimator \hat{x} , define $\hat{T} := \{j : |\hat{x}_j| \ge \mu/2\}$ Whenever $j \in T \setminus \hat{T}$ or $j \in \hat{T} \setminus T$, $|\hat{x}_j - x_j| \ge \mu/2$

$$\|\widehat{x} - x\|_{2}^{2} \ge \frac{\mu^{2}}{4} |T \setminus \widehat{T}| + \frac{\mu^{2}}{4} |\widehat{T} \setminus T| = \frac{\mu^{2}}{4} |\widehat{T} \Delta T|$$

$$\mathbb{E} \|\widehat{x} - x\|_2^2 \ge \frac{\mu^2}{4} \mathbb{E} |\widehat{T} \Delta T|$$

Proof of Main Result

Lemma Under the Bernoulli prior, any estimate \widehat{T} satisfies $\mathbb{E}\left|\widehat{T}\Delta T\right| \ge S\left(1 - \frac{\mu}{2}\sqrt{\frac{M}{N}}\right).$ Thus, $\mathbb{E} \|\widehat{x} - x\|_2^2 \ge \frac{\mu^2}{4} \mathbb{E} |\widehat{T} \Delta T|$ $\geq S \cdot \frac{\mu^2}{4} \left(1 - \frac{\mu}{2} \sqrt{\frac{M}{N}} \right)$ Plug in $\mu = \frac{8}{3} \sqrt{\frac{N}{M}}$ and this reduces to $\mathbb{E} \|\widehat{x} - x\|_{2}^{2} \ge \frac{4}{27} \cdot \frac{SN}{M} \ge \frac{1}{7} \cdot \frac{SN}{M}$

Key Ideas in Proof of Lemma

$$\mathbb{P}_{0,j}(y_1,\ldots,y_m) = \mathbb{P}(y_1,\ldots,y_m | x_j = 0)$$
$$\mathbb{P}_{1,j}(y_1,\ldots,y_m) = \mathbb{P}(y_1,\ldots,y_m | x_j = \mu)$$

$$\mathbb{E} |\widehat{T}\Delta T| \geq \frac{S}{N} \sum_{j} (1 - \|\mathbb{P}_{1,j} - \mathbb{P}_{0,j}\|_{\mathrm{TV}})$$
$$\geq S - \frac{S}{\sqrt{N}} \sqrt{\sum_{j} \|\mathbb{P}_{1,j} - \mathbb{P}_{0,j}\|_{\mathrm{TV}}^2}$$

$$\sum_{j} \|\mathbb{P}_{1,j} - \mathbb{P}_{0,j}\|_{\mathrm{TV}}^2 \le \frac{\mu^2}{4} M \longrightarrow \mathbb{E} |\widehat{T}\Delta T| \ge S \left(1 - \frac{\mu}{2}\sqrt{\frac{M}{N}}\right)$$

Key Ideas in Proof of Lemma

Pinsker's Inequality

$$\|\mathbb{P} - \mathbb{Q}\|_{\mathrm{TV}} \le \sqrt{K(\mathbb{P}, \mathbb{Q})/2}$$

$$\begin{aligned} \|\mathbb{P}_{1,j} - \mathbb{P}_{0,j}\|_{\mathrm{TV}}^2 &\leq \frac{\pi_0}{2} K(\mathbb{P}_{0,j}, \mathbb{P}_{1,j}) + \frac{\pi_1}{2} K(\mathbb{P}_{1,j}, \mathbb{P}_{0,j}) \\ &\leq \frac{\mu^2}{4} \sum_i \mathbb{E} \phi_{i,j}^2 \end{aligned}$$

$$\sum_{j} \|\mathbb{P}_{1,j} - \mathbb{P}_{0,j}\|_{\mathrm{TV}}^2 \le \frac{\mu^2}{4} \sum_{i,j} \mathbb{E} \phi_{i,j}^2 = \frac{\mu^2}{4} M$$

Suppose that S = 1 and that $x_{j^*} = \mu$

- split measurements into $\log N$ stages
- in each stage, use measurements to decide if the nonzero is in the left or right half of the "active set"
- after subdividing $\log N$ times, return support



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Experimental Results



[Arias-Castro, Candès, and Davenport - 2011]

Open Questions

- No method can succeed when $\frac{\mu}{\sigma} \approx \sqrt{N/M}$, but the binary search approach succeeds as long as $\frac{\mu}{\sigma} \geq C\sqrt{N/M}$ [Davenport and Arias-Castro; Malloy and Nowak - 2012]
- Practical algorithms that work well for all values of $\boldsymbol{\mu}$
- Optimal algorithms for S>1
- New theory for restricted adaptive measurements
 - single-pixel camera: 0/1 measurements
 - magnetic resonance imaging (MRI): Fourier measurements
 - analog-to-digital converters: linear filter measurements
- New sensors and architectures that can actually acquire adaptive measurements

More Information

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