## Streaming solutions to least-squares problems

In our discussion of least-squares so far, we have focussed on static problems: a set of measurements $\boldsymbol{y}=\boldsymbol{A} \boldsymbol{x}_{0}+\boldsymbol{e}$ comes in all at once, and we use them all to estimate $\boldsymbol{x}_{0}$.

In this section, we will shift our focus to streaming problems. We observe

$$
\begin{aligned}
\boldsymbol{y}_{0} & =\boldsymbol{A}_{0} \boldsymbol{x}_{1}+\boldsymbol{e}_{0} \\
\boldsymbol{y}_{1} & =\boldsymbol{A}_{1} \boldsymbol{x}_{2}+\boldsymbol{e}_{1} \\
\vdots & \\
\boldsymbol{y}_{k} & =\boldsymbol{A}_{k} \boldsymbol{x}_{k}+\boldsymbol{e}_{k}
\end{aligned}
$$

At each time $k$, we want to form the best estimate of $\boldsymbol{x}_{k}$ from the observations $\boldsymbol{y}_{0}, \boldsymbol{y}_{1}, \ldots, \boldsymbol{y}_{k}$ seen up to that point exploiting some known structure regarding how the $\boldsymbol{x}_{k}$ are related to each other. Moreover, we would like to do this in an efficient manner. The size of the problem is growing with $k$ - rather than resolving the problem from scratch every time, we would like a principled (and fast) way to update the solution when a new observation is made.

We will consider two basic frameworks:

1. Recursive Least Squares (RLS):

The vector $\boldsymbol{x}_{k}=\boldsymbol{x}_{*}$ for all $k$, i.e. we are estimating a static vector $\boldsymbol{x}_{*}$.
2. The Kalman filter:

The vector $\boldsymbol{x}_{k}$ moves at every times step, and we have a (linear) dynamical model for how it moves.

In both of these frameworks, the measurements matrices $\boldsymbol{A}_{1}, \boldsymbol{A}_{2}, \ldots$ can be different, and can even have a different number of rows. We will assume that the total number of measurements we have seen at any point exceeds the number of unknowns, and if we form

$$
\underline{\boldsymbol{A}}_{k}=\left[\begin{array}{c}
\boldsymbol{A}_{0} \\
\boldsymbol{A}_{1} \\
\vdots \\
\boldsymbol{A}_{k}
\end{array}\right]
$$

then $\underline{\boldsymbol{A}}_{k}^{\mathrm{T}} \underline{\boldsymbol{A}}_{k}$ is invertible. Generalizing what we say to rank-deficient systems is not hard, but this assumption makes the discussion easier.

The key piece of mathematical technology we need is the matrix inversion lemma.

## The Matrix Inversion Lemma

The matrix inversion lemma shows us how the solution to a system of equations can be efficiently updated. Here we state a slightly simplified version of the result: If $\boldsymbol{W}$ is an $N \times N$ invertible matrix and $\boldsymbol{X}$ is an $R \times N$ matrix, then the following identity holds:

$$
\left(\boldsymbol{W}+\boldsymbol{X}^{\mathrm{T}} \boldsymbol{X}\right)^{-1}=\boldsymbol{W}^{-1}-\boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\left(\mathbf{I}+\boldsymbol{X} \boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\right)^{-1} \boldsymbol{X} \boldsymbol{W}^{-1}
$$

This is a special case of the Sherman-Morrison-Woodbury identity, and is straightforward to prove (see the Technical Details at the end of these notes). The point is that if $\boldsymbol{W}^{-1}$ has already been calculated, then finding a solution to $\left(\boldsymbol{W}+\boldsymbol{X}^{\mathrm{T}} \boldsymbol{X}\right) \boldsymbol{w}=\boldsymbol{v}$ costs $O\left(N^{2} R\right)+O\left(N R^{2}\right)+O\left(R^{3}\right)$ instead of $O\left(N^{3}\right)$. If $R$ is very small compared to $N$, this can be a significant savings.

## Updating least-squares solutions

We can apply this fact to efficiently update the solution to leastsquares problems as new measurements become available.

Suppose we have observed

$$
\boldsymbol{y}_{0}=\boldsymbol{A}_{0} \boldsymbol{x}_{*}+\boldsymbol{e}_{0}
$$

and have formed the least-squares estimate

$$
\widehat{\boldsymbol{x}}_{0}=\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}\right)^{-1} \boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{y}_{0} .
$$

Now we observe

$$
\boldsymbol{y}_{1}=\boldsymbol{A}_{1} \boldsymbol{x}_{*}+\boldsymbol{e}_{1},
$$

where $\boldsymbol{A}_{1}$ is an $M_{1} \times N$ matrix with $M_{1} \ll N$. Given $\boldsymbol{y}_{0}$ and $\boldsymbol{y}_{1}$, the full least-squares estimate is formed from the system of equations

$$
\left[\begin{array}{l}
\boldsymbol{y}_{0} \\
\boldsymbol{y}_{1}
\end{array}\right]=\left[\begin{array}{l}
\boldsymbol{A}_{0} \\
\boldsymbol{A}_{1}
\end{array}\right] \boldsymbol{x}_{*}+\left[\begin{array}{l}
\boldsymbol{e}_{0} \\
\boldsymbol{e}_{1}
\end{array}\right],
$$

resulting in

$$
\widehat{\boldsymbol{x}}_{1}=\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{A}_{1}\right)^{-1}\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{y}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{y}_{1}\right) .
$$

Now let $\boldsymbol{P}_{k}$ be the inverse of the aggregated system we would like to solve at each step:

$$
\begin{aligned}
& \boldsymbol{P}_{0}=\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}\right)^{-1} \\
& \boldsymbol{P}_{1}=\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{A}_{1}\right)^{-1} .
\end{aligned}
$$

Then using the matrix inversion lemma with $\boldsymbol{W}=\boldsymbol{A}_{0}^{T} \boldsymbol{A}_{0}=\boldsymbol{P}_{0}^{-1}$ and $\boldsymbol{X}=\boldsymbol{A}_{1}$ gives us the update

$$
\boldsymbol{P}_{1}=\boldsymbol{P}_{0}-\boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}\left(\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}\right)^{-1} \boldsymbol{A}_{1} \boldsymbol{P}_{0} .
$$

When the number of new measurements (rows in $\boldsymbol{A}_{1}$ ) is small, then the system of equations $\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}$ can be much easier to handle than $\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{A}_{1}=\boldsymbol{P}_{1}^{-1}$. For example, suppose we see just one new measurement, so the matrix $\boldsymbol{A}_{1}$ has just one row: $\boldsymbol{A}_{1}=\boldsymbol{a}_{1}^{\mathrm{T}}$, $\boldsymbol{a}_{1} \in \mathbb{R}^{N}$. Then

$$
y_{1}=\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{x}_{*}+e_{1},
$$

and

$$
\widehat{\boldsymbol{x}}_{1}=\left[\boldsymbol{P}_{0}-\boldsymbol{P}_{0} \boldsymbol{a}_{1}\left(1+\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{P}_{0} \boldsymbol{a}_{1}\right)^{-1} \boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{P}_{0}\right]\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{y}_{0}+y_{1} \boldsymbol{a}_{1}\right) .
$$

Set $\boldsymbol{u}=\boldsymbol{P}_{0} \boldsymbol{a}_{1}$. Then

$$
\begin{aligned}
\widehat{\boldsymbol{x}}_{1} & =\widehat{\boldsymbol{x}}_{0}+y_{1} \boldsymbol{u}-\frac{\boldsymbol{a}_{1}^{\mathrm{T}} \widehat{\boldsymbol{x}}_{0}}{1+\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{u}} \boldsymbol{u}-\frac{y_{1} \cdot \boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{u}}{1+\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{u}} \boldsymbol{u} \\
& =\widehat{\boldsymbol{x}}_{0}+\left(\frac{1}{1+\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{u}}\right)\left(y_{1}-\boldsymbol{a}_{1}^{\mathrm{T}} \widehat{\boldsymbol{x}}_{0}\right) \boldsymbol{u} .
\end{aligned}
$$

Thus we can update the solution with one vector-matrix multiply (which has cost $O\left(N^{2}\right)$ ) and two inner products (with cost $O(N)$ ).

In addition, we can carry forward the "information matrix" using the update

$$
\boldsymbol{P}_{1}=\boldsymbol{P}_{0}-\frac{1}{1+\boldsymbol{a}_{1}^{\mathrm{T}} \boldsymbol{u}} \boldsymbol{u} \boldsymbol{u}^{\mathrm{T}}
$$

In general (for $M_{1}$ new measurements), we have

$$
\begin{aligned}
\widehat{\boldsymbol{x}}_{1} & =\boldsymbol{P}_{1}\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{y}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{y}_{1}\right) \\
& =\boldsymbol{P}_{1}\left(\boldsymbol{P}_{0}^{-1} \widehat{\boldsymbol{x}}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{y}_{1}\right),
\end{aligned}
$$

and since

$$
\boldsymbol{P}_{0}^{-1}=\boldsymbol{P}_{1}^{-1}-\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{A}_{1}
$$

this implies

$$
\begin{aligned}
\widehat{\boldsymbol{x}}_{1} & =\boldsymbol{P}_{1}\left(\boldsymbol{P}_{1}^{-1} \widehat{\boldsymbol{x}}_{0}-\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{A}_{1} \widehat{\boldsymbol{x}}_{0}+\boldsymbol{A}_{1}^{\mathrm{T}} \boldsymbol{y}_{1}\right) \\
& =\widehat{\boldsymbol{x}}_{0}+\boldsymbol{K}_{1}\left(\boldsymbol{y}_{1}-\boldsymbol{A}_{1} \widehat{\boldsymbol{x}}_{0}\right),
\end{aligned}
$$

where $\boldsymbol{K}_{1}$ is the "gain matrix"

$$
\boldsymbol{K}_{1}=\boldsymbol{P}_{1} \boldsymbol{A}_{1}^{\mathrm{T}}
$$

The update for $\boldsymbol{P}_{1}$ is

$$
\begin{aligned}
\boldsymbol{P}_{1} & =\boldsymbol{P}_{0}-\boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}\left(\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}\right)^{-1} \boldsymbol{A}_{1} \boldsymbol{P}_{0} \\
& =\boldsymbol{P}_{0}-\boldsymbol{U}\left(\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{U}\right)^{-1} \boldsymbol{U}^{\mathrm{T}},
\end{aligned}
$$

where $\boldsymbol{U}=\boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}$ is an $N \times M_{1}$ matrix, and $\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{U}$ is $M_{1} \times M_{1}$. So the cost of the update is

- $O\left(M_{1} N^{2}\right)$ to compute $\boldsymbol{U}=\boldsymbol{P}_{0} \boldsymbol{A}_{1}^{\mathrm{T}}$,
- $O\left(M_{1}^{2} N\right)$ to compute $\boldsymbol{A}_{1} \boldsymbol{U}$,
- $O\left(M_{1}^{3}\right)$ to invert ${ }^{1}\left(\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{U}\right)^{-1}$,
- $O\left(M_{1}^{2} N\right)$ to compute $\left(\mathbf{I}+\boldsymbol{A}_{1} \boldsymbol{U}\right)^{-1} \boldsymbol{U}^{\mathrm{T}}$,
- $O\left(M_{1} N^{2}\right)$ to take the result of the last step and apply $\boldsymbol{U}$,
- $O\left(N^{2}\right)$ to subtract the result of the last step from $\boldsymbol{P}_{0}$.

So assuming that $M_{1}<N$, the overall cost is $O\left(M_{1} N^{2}\right)$, which is on the order of $M_{1}$ vector-matrix multiplies.

[^0]
## Recursive Least Squares (RLS)

## Given

$$
\begin{aligned}
& \boldsymbol{y}_{0}=\boldsymbol{A}_{0} \boldsymbol{x}_{*}+\boldsymbol{e}_{0} \\
& \boldsymbol{y}_{1}=\boldsymbol{A}_{1} \boldsymbol{x}_{*}+\boldsymbol{e}_{1} \\
& \vdots \\
& \boldsymbol{y}_{k}=\boldsymbol{A}_{k} \boldsymbol{x}_{*}+\boldsymbol{e}_{k} \\
& \vdots
\end{aligned}
$$

RLS is an online algorithm for computing the best estimate for $\boldsymbol{x}_{*}$ from all the measurements it has seen up to the current time.

$$
\begin{aligned}
& \text { Recursive Least Squares } \\
& \text { Initialize: ( } \boldsymbol{y}_{0} \text { appears) } \\
& \boldsymbol{P}_{0}=\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{A}_{0}\right)^{-1} \\
& \widehat{\boldsymbol{x}}_{0}=\boldsymbol{P}_{0}\left(\boldsymbol{A}_{0}^{\mathrm{T}} \boldsymbol{y}_{0}\right) \\
& \text { for } k=1,2,3, \ldots \text { do } \\
& \text { ( } \boldsymbol{y}_{k} \text { appears) } \\
& \boldsymbol{P}_{k}=\boldsymbol{P}_{k-1}-\boldsymbol{P}_{k-1} \boldsymbol{A}_{k}^{\mathrm{T}}\left(\mathbf{I}+\boldsymbol{A}_{k} \boldsymbol{P}_{k-1} \boldsymbol{A}_{k}^{\mathrm{T}}\right)^{-1} \boldsymbol{A}_{k} \boldsymbol{P}_{k-1} \\
& \boldsymbol{K}_{k}=\boldsymbol{P}_{k} \boldsymbol{A}_{k}^{\mathrm{T}} \\
& \widehat{\boldsymbol{x}}_{k}=\widehat{\boldsymbol{x}}_{k-1}+\boldsymbol{K}_{k}\left(\boldsymbol{y}_{k}-\boldsymbol{A}_{k} \widehat{\boldsymbol{x}}_{k-1}\right) \\
& \text { end for }
\end{aligned}
$$

## Technical Details: Matrix Inversion Lemma

The general statement of the Sherman-Morrison-Woodbury identity is that

$$
\left(\boldsymbol{W}+\boldsymbol{X}^{\mathrm{T}} \boldsymbol{Y} \boldsymbol{Z}\right)^{-1}=\boldsymbol{W}^{-1}-\boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\left(\boldsymbol{Y}^{-1}+\boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\right)^{-1} \boldsymbol{Z} \boldsymbol{W}^{-1}
$$

where $\boldsymbol{W}$ is $N \times N$ and invertible, $\boldsymbol{X}$ and $\boldsymbol{Z}$ are $R \times N$, and $\boldsymbol{Y}$ is $R \times R$ and invertible.

The proof of this is straightforward. Given any right hand side $\boldsymbol{v} \in$ $\mathbb{R}^{N}$, we would like to solve

$$
\begin{equation*}
\left(\boldsymbol{W}+\boldsymbol{X}^{\mathrm{T}} \boldsymbol{Y} \boldsymbol{Z}\right) \boldsymbol{w}=\boldsymbol{v} \tag{1}
\end{equation*}
$$

for $\boldsymbol{w}$. Set

$$
\boldsymbol{z}=\boldsymbol{Y} \boldsymbol{Z} \boldsymbol{w} \quad \Rightarrow \quad \boldsymbol{Y}^{-1} \boldsymbol{z}=\boldsymbol{Z} \boldsymbol{w}
$$

We now have the set of two equations

$$
\begin{aligned}
\boldsymbol{W} \boldsymbol{w}+\boldsymbol{X}^{\mathrm{T}} \boldsymbol{z} & =\boldsymbol{v} \\
\boldsymbol{Z} \boldsymbol{w}-\boldsymbol{Y}^{-1} \boldsymbol{z} & =\mathbf{0}
\end{aligned}
$$

Manipulating the first equation yields

$$
\begin{equation*}
\boldsymbol{w}=\boldsymbol{W}^{-1}\left(\boldsymbol{v}-\boldsymbol{X}^{\mathrm{T}} \boldsymbol{z}\right) \tag{2}
\end{equation*}
$$

and then plugging this into the second equation gives us $\boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{v}-\boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}} \boldsymbol{z}-\boldsymbol{Y}^{-1} \boldsymbol{z}=\mathbf{0}$

$$
\begin{equation*}
\Rightarrow \quad \boldsymbol{z}=\left(\boldsymbol{Y}^{-1}+\boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\right)^{-1} \boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{v} \tag{3}
\end{equation*}
$$

So then given any $\boldsymbol{v} \in \mathbb{R}^{N}$, we can solve for $\boldsymbol{w}$ in (1) by combining (2) and (3) to get

$$
\boldsymbol{w}=\boldsymbol{W}^{-1} \boldsymbol{v}-\boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\left(\boldsymbol{Y}^{-1}+\boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{X}^{\mathrm{T}}\right)^{-1} \boldsymbol{Z} \boldsymbol{W}^{-1} \boldsymbol{v}
$$

As this holds for any right-hand side $\boldsymbol{v}$, this establishes the result.


[^0]:    ${ }^{1}$ In practice, it is probably more stable to find and update a factorization of this matrix. But the cost is the same.

