# How well can we estimate a sparse vector?

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### **Sparse Estimation**



How well can we estimate x ?

# **Applications**

- Statistics
  - model selection / variable selection in high-dimensional regression

• Inverse problems

- Compressive sensing (CS)
  - matrix  $\Phi$  represents a sensing system
  - typically underdetermined
  - sparsity acts as a regularizer

# Core Challenges in CS

• How should we design the matrix  $\Phi$  so that M is as small as possible?



• How can we recover x from the measurements y?

#### Answers

- Choose a random matrix
  - fill out the entries of  $\Phi$  with i.i.d. samples from a sub-Gaussian distribution
  - project onto a "random subspace"

$$M = O(S \log(N/S)) \ll N$$

• Use any sparse approximation algorithm

Is this the best we can do?

# **Recovery from Noisy Measurements**

Given 
$$y = \Phi x + e$$
 or  $y = \Phi(x + n)$ ,  
find  $x$ 

- Optimization-based methods
  - basis pursuit, basis pursuit de-noising, Dantzig selector

$$\widehat{x} = \underset{x \in \mathbb{R}^{N}}{\arg\min} \|x\|_{1}$$
  
s.t. 
$$\|y - \Phi x\|_{2} \le \epsilon$$

- Greedy/Iterative algorithms
  - OMP, StOMP, ROMP, CoSaMP, Thresh, SP, IHT, ...

## Stable Signal Recovery

Suppose that we observe  $y = \Phi x + e$  and that  $\Phi$  satisfies the RIP of order 2S.

$$(1-\delta)\|x\|_2^2 \le \|\Phi x\|_2^2 \le (1+\delta)\|x\|_2^2 \qquad \|x\|_0 \le 2S$$

Typical (worst-case) guarantee

$$\|\widehat{x} - x\|_2^2 \le C \|e\|_2^2$$

Even if  $\Lambda = \operatorname{supp}(x)$  is provided by an oracle, the error can still be as large as  $\|\widehat{x} - x\|_2^2 = \|e\|_2^2/(1-\delta)$ .

# Stable Signal Recovery: Part II

Suppose now that  $\Phi$  satisfies

$$A(1-\delta)\|x\|_{2}^{2} \leq \|\Phi x\|_{2}^{2} \leq A(1+\delta)\|x\|_{2}^{2} \qquad \|x\|_{0} \leq 2S$$

#### In this case our guarantee becomes

$$\|\widehat{x} - x\|_2^2 \le \frac{C}{A} \|e\|_2^2$$

Unit-norm rows  $\|\widehat{x} - x\|_2^2 \leq C \frac{N}{M} \|e\|_2^2$ 

# **Expected Performance**

- Worst-case bounds can be pessimistic
- What about the *average* error?
  - assume e is white noise with variance  $\sigma^2$

 $\mathbb{E}\left(\|e\|_2^2\right) = M\sigma^2$ 

- for (nonadaptive) oracle

$$\mathbb{E}\left(\|\widehat{x} - x\|_2^2\right) \le \frac{S\sigma^2}{A(1-\delta)}$$

- if e is Gaussian, then for  $\ell_1$  -minimization

$$\mathbb{E}\left(\|\widehat{x} - x\|_{2}^{2}\right) \leq \frac{C'}{A}S\sigma^{2}\log N$$

#### Can We Do Better?

- Better choice of  $\Phi$  ?
- Better recovery algorithm?

Assume we have a budget for  $\|\Phi\|_F^2$ .

If we knew the support of x a priori, then by adapting  $\Phi$  to exploit this knowledge we could achieve

$$\mathbb{E}\left[\|\widehat{x} - x\|_2^2\right] \approx \frac{S}{\|\Phi\|_F^2} S\sigma^2 \ll C' \frac{N}{\|\Phi\|_F^2} S\sigma^2 \log N$$

Is there any way to match this performance without knowing the support of x in advance?

$$R^*_{\mathrm{mm}}(\Phi) = \inf_{\widehat{x}} \sup_{\|x\|_0 \le S} \mathbb{E}\left[\|\widehat{x}(\Phi x + e) - x\|_2^2\right]$$

#### No!

Theorem:  
If 
$$y = \Phi x + e$$
 with  $e \sim \mathcal{N}(0, \sigma^2 I)$ , then  
 $R_{mm}^*(\Phi) \ge C \frac{N}{\|\Phi\|_F^2} S \sigma^2 \log(N/S)$ .  
If  $y = \Phi(x+n)$  with  $n \sim \mathcal{N}(0, \sigma^2 I)$ , then  
 $R_{mm}^*(\Phi) \ge C \frac{N}{M} S \sigma^2 \log(N/S)$ .

$$\Phi = U\Sigma V^* \quad y' = \Sigma^{-1} U^* y = V^* x + V^* n \quad \|V^*\|_F^2 = M$$

See also: Raskutti, Wainwright, and Yu (2009) Ye and Zhang (2010)

[Candès and Davenport - 2011]

# Intuition

Suppose that y = x + e with  $e \sim \mathcal{N}(0, I)$  and that S = 1.

 $R^*_{\rm mm}(I) \ge C \log(N).$ 



# **Proof Recipe**

Ingredients [Makes  $\sigma^2 = 1$  servings]

- Lemma 1: Suppose  $\mathcal{X}$  is a set of S-sparse points such that  $\|x_i x_j\|_2^2 \ge 8NR_{\min}^*(\Phi)$  for all  $x_i, x_j \in \mathcal{X}$ . Then  $\frac{1}{2} \log |\mathcal{X}| - 1 \le \frac{1}{2|\mathcal{X}|^2} \sum_{i,j} \|\Phi x_i - \Phi x_j\|_2^2$ .
- Lemma 2: There exists a set  $\mathcal{X}$  of S-sparse points such that

$$\begin{aligned} \bullet & |\mathcal{X}| = (N/S)^{S/4} \\ \bullet & \|x_i - x_j\|_2 \geq \frac{1}{2} \text{ for all } x_i, x_j \in \mathcal{X} \\ \bullet & \|\frac{1}{|\mathcal{X}|} \sum_i x_i x_i^* - \frac{1}{N}I\| \leq \frac{\beta}{N} \text{ for some } \beta > 0 \end{aligned}$$

#### Instructions

Combine ingredients and add a dash of linear algebra.

#### **Proof Outline**

$$\mu = \frac{1}{|\mathcal{X}|} \sum_{i} x_{i} \quad Q = \frac{1}{|\mathcal{X}|} \sum_{i} x_{i} x_{i}^{*}$$

 $\frac{S}{4} \log(N/S) - 2 \le \frac{1}{|\mathcal{X}|^2} \sum_{i,j} \|\Phi x_i - \Phi x_j\|_2^2$  $= \operatorname{Tr}\left(\Phi^*\Phi\left(\frac{1}{|\mathcal{X}|^2}\sum_{i,j}(x_i - x_j)(x_i - x_j)^*\right)\right)$  $= \operatorname{Tr} \left( \Phi^* \Phi \left( 2(Q - \mu \mu^*) \right) \right)$  $\leq 2 \operatorname{Tr} (\Phi^* \Phi Q)$  $\leq 2 \operatorname{Tr} (\Phi^* \Phi) \| Q \|$  $< 2 \|\Phi\|_{F}^{2} \cdot 16 R_{mm}^{*}(\Phi)(1+\beta)$  $R_{\rm mm}^*(\Phi) \ge \frac{S \log(N/S)}{128(1+\beta) \|\Phi\|^2}$ 

## Recall: Lemma 2

Lemma 2: There exists a set  $\mathcal{X}$  of S-sparse points such that

• 
$$|\mathcal{X}| = (N/S)^{S/4}$$
  
•  $||x_i - x_j||_2 \ge \frac{1}{2}$  for all  $x_i, x_j \in \mathcal{X}$   
•  $||\frac{1}{|\mathcal{X}|} \sum_i x_i x_i^* - \frac{1}{N}I|| \le \frac{\beta}{N}$  for some  $\beta > 0$ 

#### Strategy

Construct  ${\mathcal X}$  by sampling (with replacement) from

$$\mathcal{U} = \left\{ x \in \{0, \sqrt{1/S}, -\sqrt{1/S}\}^N : \|x\|_0 \le S \right\}$$

Repeat for  $|\mathcal{X}| = (N/S)^{S/4}$  iterations.

With probability > 0, the remaining properties are satisfied.

Key: Matrix Bernstein Inequality [Ahlswede and Winter, 2002]

#### Recap

Noise added to the *measurements* 

$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \leq C' \frac{N}{\|\Phi\|_{F}^{2}} S\sigma^{2} \log N$$
$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \geq C \frac{N}{\|\Phi\|_{F}^{2}} S\sigma^{2} \log(N/S)$$

Noise added to the *signal* 

$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \leq C' \frac{N}{M} S \sigma^{2} \log N$$
$$\mathbb{E}\left[\|\widehat{x} - x\|_{2}^{2}\right] \geq C \frac{N}{M} S \sigma^{2} \log(N/S)$$

# **Noise Folding**



[Davenport, Laska, Treichler, and Baraniuk - 2011]

# Adaptivity to the Rescue?

What if we adapt the measurements to the particular signal?



If we are too greedy, our support estimate might be wrong...

Does adaptivity really help?

## Sometimes...

- Information-based complexity: "Adaptivity doesn't help!"
  - assumes signal x lies in a set K satisfying certain conditions
  - noise-free measurements
  - adaptivity reduces minimax error over  $\,K\,$  by at most  $2\,$
- Nevertheless, adaptivity can still help [Indyk et al. 2011]
  - reduced number of measurements in a probabilistic setting
  - still requires noise-free measurements
- What about noise?
  - distilled sensing (Haupt, Castro, Nowak, and others)
  - message seems to be that adaptivity really helps in noise

# Adaptive Compressive Sensing

Suppose we have a budget of M measurements of the form

$$y_i = \langle \phi_i, x \rangle + e_i$$

where  $\|\phi_i\|_2 = 1$  and  $e_i \sim \mathcal{N}(0, \sigma^2)$ .

The vector  $\phi_i$  can have an arbitrary (but deterministic) dependence on the measurements  $y_1, y_2, \ldots, y_{i-1}$ .

Consider the minimax MSE

$$R_{\mathrm{mm}}^* = \inf_{\widehat{x}} \sup_{\|x\|_0 \le S} \mathbb{E} \left[ \|\widehat{x}(\Phi x + e) - x\|_2^2 \right]$$

## Main Result

Possibilities include

- Adaptive oracle rate:  $R_{\rm mm}^* \approx \frac{S}{M}S\sigma^2$
- Nonadaptive rate:  $R^*_{\rm mm} \approx \frac{N}{M} S \sigma^2 \log(N/S)$
- Somewhere in-between?

$$R^*_{\rm mm} \ge \tilde{C} \frac{N}{M} S \sigma^2$$

In general, adaptivity does *not* significantly help!!

# **Underlying Ideas**

- **Step 1:** Consider sparse signals with nonzeros of amplitude  $\mu = \sqrt{N/M}$ .
- Step 2: Show that if you have fewer than M measurements, then with high probability you will fail to recover a significant fraction of the support.
- Step 3: Immediately translate this into a lower bound on the MSE.

# Adaptivity in Practice

Suppose that S = 1 and that  $x_{j^*} = \mu$ .

Algorithm 1 [Castro et al. - 2008]

- start with random (Rademacher) measurements
- after each measurement, compute posterior distribution p
- re-weight subsequent measurements using p

Algorithm 2 [Iwen and Tewfik - 2011]

- split measurements into  $\log N$  stages
- in each stage, use measurements to decide if the nonzero is in the left or right half of the "active set"
- after subdividing  $\log N$  times, return support

#### Phase Transition in the Posterior



 $n = 512 \quad \sigma^2 = 1$ 

#### Phase Transition in the MSE



# Conclusions

- In some scenarios, CS can be sensitive to noise
  - inherent lower bound that applies to any possible sensing scheme
  - if you can average out noise, that will always help
  - sparsity is still helping a lot
- Surprisingly, adaptive algorithms cannot overcome this obstacle!
- Adaptivity might still be very useful in practice
  - practical adaptive algorithms that achieve the minimax rate for all values of  $\,\mu$  ?