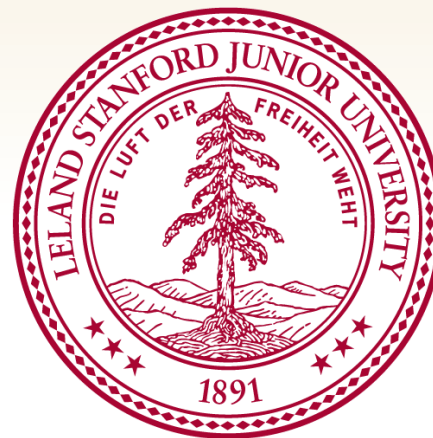


# Compressive Sensing

## Part III: Compressive Sensing in Practice

*Mark A. Davenport*

Stanford University  
Department of Statistics

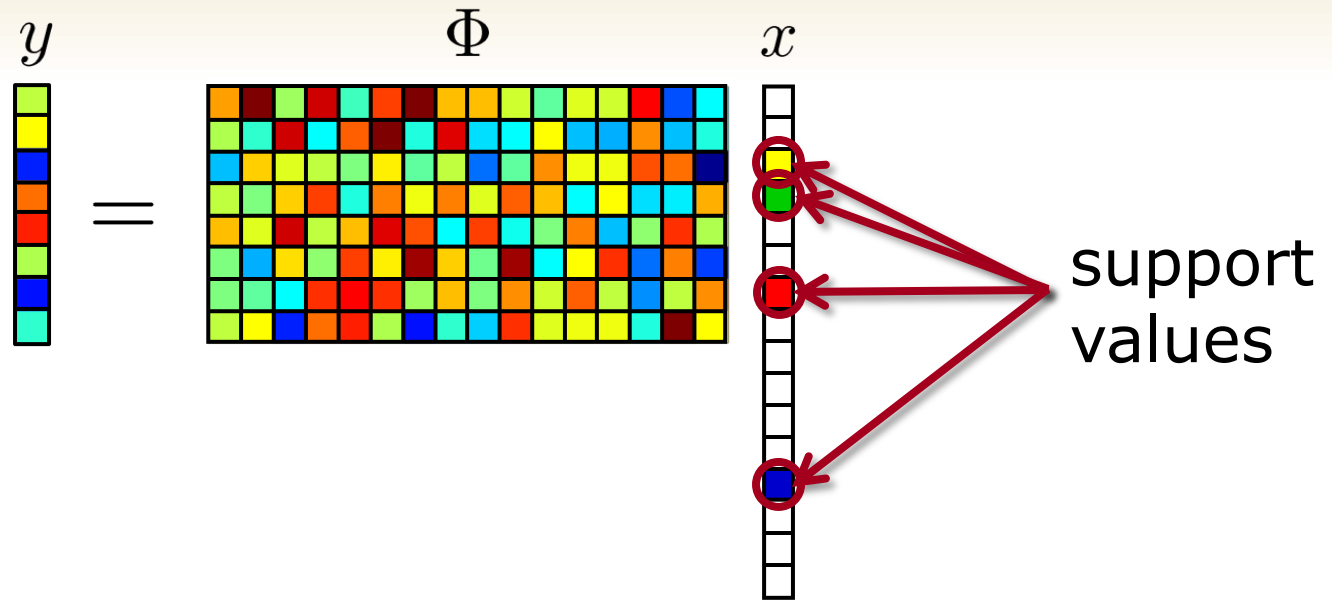


# Important Practical Challenges

- Noise!
  - noisy measurements
  - noisy signals
  - interference
- Quantization
  - quantization error
  - saturation effects
- Good signal models
  - is sparsity sometimes not enough?
  - what dictionaries should we use in practice?

# **Measurement and Signal Noise**

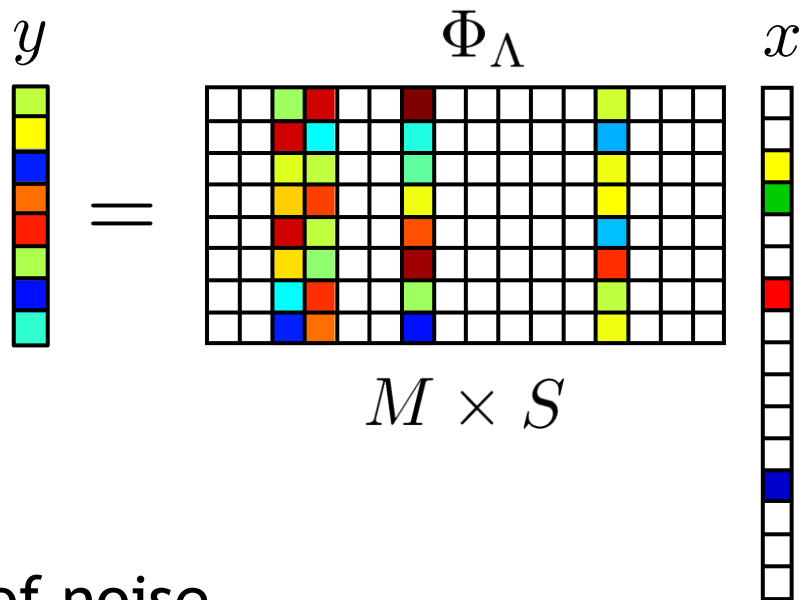
# Sparse Signal Recovery



- Optimization /  $\ell_1$  -minimization
- Greedy algorithms
  - matching pursuit
  - orthogonal matching pursuit (OMP)
  - regularized OMP
  - CoSaMP, Subspace Pursuit, IHT, ...

# Exact Recovery

If we can determine  $\Lambda = \text{supp}(x)$ , then the problem becomes *over-determined*.



In the absence of noise,

$$\begin{aligned}\Phi_\Lambda^\dagger y &= (\Phi_\Lambda^T \Phi_\Lambda)^{-1} \Phi_\Lambda^T y \\ &= (\Phi_\Lambda^T \Phi_\Lambda)^{-1} \Phi_\Lambda^T \Phi_\Lambda x \\ &= x\end{aligned}$$

# Signal Recovery in Noise

Given  $y = \Phi x + e$   
find  $x$

- Optimization-based methods
  - basis pursuit, basis pursuit de-noising, Dantzig selector

$$\hat{x} = \arg \min_{x \in \mathbb{R}^N} \|x\|_1$$
$$\text{s.t. } \|y - \Phi x\|_2 \leq \epsilon$$

- Greedy/Iterative algorithms
  - OMP, StOMP, ROMP, CoSaMP, Thresh, SP, IHT, ...

# Stable Signal Recovery

Suppose that we observe  $y = \Phi x + e$  and that  $\Phi$  satisfies the RIP of order  $S$ .

Typical (worst-case) guarantee

$$\|\hat{x} - x\|_2 \leq C\|e\|_2$$

Even if  $\Lambda = \text{supp}(x)$  is provided by an oracle, the error can still be as large as

$$\|\hat{x} - x\|_2 = \frac{\|e\|_2}{1 - \delta}$$

# Expected Performance

- Worst-case bounds can be pessimistic
- What about the *average* error?
  - assume  $e$  is white noise with variance  $\sigma^2$

$$\mathbb{E} (\|e\|_2^2) = M\sigma^2$$

- for oracle-assisted estimator

$$\mathbb{E} (\|\hat{x} - x\|_2) \leq \frac{S\sigma^2}{1 - \delta}$$

- if  $e$  is Gaussian, then for  $\ell_1$ -minimization

$$\mathbb{E} (\|\hat{x} - x\|_2^2) \leq CS\sigma^2 \log N$$



# White Signal Noise

What if our signal  $x$  is contaminated with noise?

$$y = \Phi(x + n)$$

Suppose  $\Phi$  satisfies the RIP and has orthogonal and equal-norm rows. If  $n$  is white noise with variance  $\sigma^2$ , then  $\Phi n$  is white noise with variance  $\sigma^2 \frac{N}{M}$ .

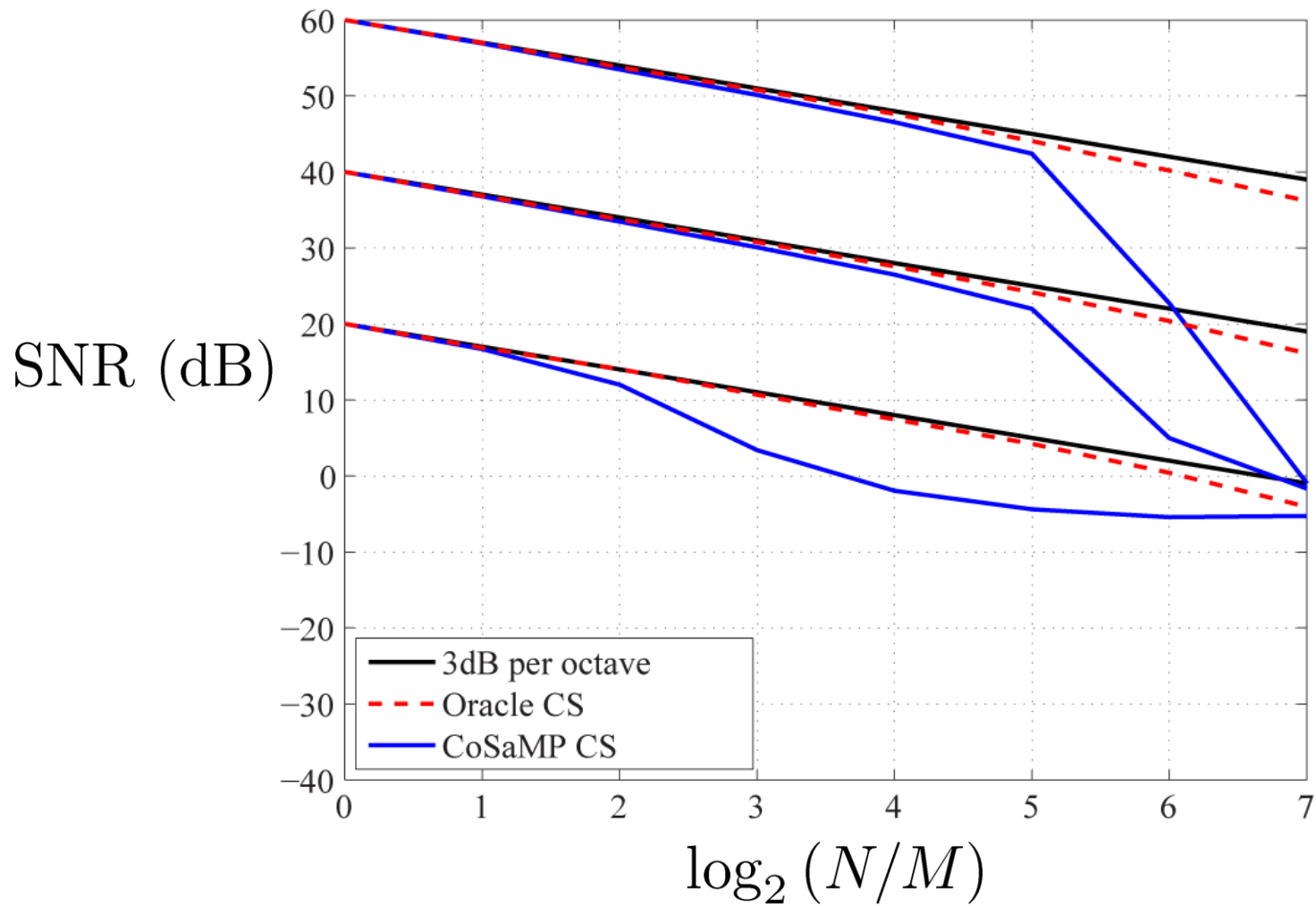
$$\|\hat{x} - x\|_2^2 \leq C' \frac{N}{M} S \sigma^2 \log N$$

$$\text{SNR} = 10 \log_{10} \left( \frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right)$$



3dB loss per octave  
of subsampling

# Noise Folding



# Can We Do Better?

- Better choice of  $\Phi$  ?
- Better recovery algorithm?

If we knew the support of  $x$  *a priori*, then we could achieve

$$\|\hat{x} - x\|_2^2 \approx \frac{S}{M} S \sigma^2 \ll C' \frac{N}{M} S \sigma^2 \log N$$

Is there any way to match this performance without knowing the support of  $x$  in advance?

$$R_{\text{mm}}^*(\Phi) = \inf_{\hat{x}} \sup_{\|x\|_0 \leq S} \mathbb{E} [\|\hat{x}(y) - x\|_2^2]$$

# No!

## Theorem:

If  $y = \Phi x + e$  with  $e \sim \mathcal{N}(0, \sigma^2 I)$ , then

$$R_{\text{mm}}^*(\Phi) \geq C \frac{N}{\|\Phi\|_F^2} S \sigma^2 \log(N/S).$$

If  $y = \Phi(x + n)$  with  $n \sim \mathcal{N}(0, \sigma^2 I)$ , then

$$R_{\text{mm}}^*(\Phi) \geq C \frac{N}{M} S \sigma^2 \log(N/S).$$

Ingredients in proof:

- Fano's inequality
- Random construction of packing set of sparse points
- Matrix Bernstein inequality to bound empirical covariance matrix of packing set

# Interference

$$y = \Phi x + e$$

- What if  $e$  represents corruption or *structured noise*, rather than Gaussian noise or arbitrary perturbations?
- Structured signal noise:

$$y = \Phi x_S + \Phi x_I$$

- Structured measurement noise:

$$y = \Phi x + \Omega e$$

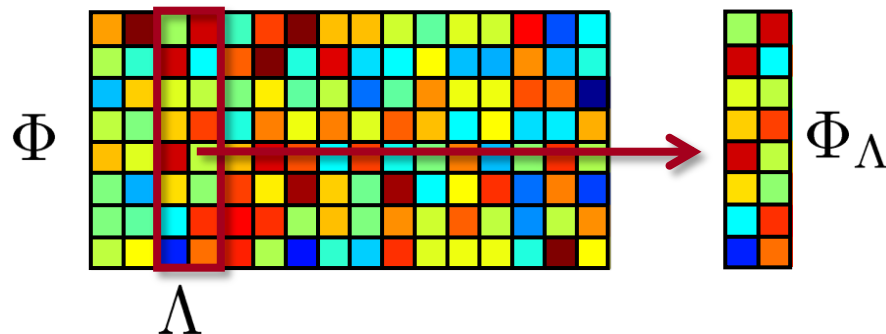
# Interference Cancellation

Suppose  $x = x_S + x_I$  where  $x_S$  is sparse with *unknown* support and  $x_I$  is sparse with *known* support  $\Lambda$

**Goal:** Design an  $M \times M$  matrix  $P$  such that

$$\|P(\Phi x_I)\|_2 \approx 0$$

$$\|P(\Phi x_S)\|_2 \approx \|\Phi x_S\|_2$$



$$P = I - \underbrace{\Phi_\Lambda \Phi_\Lambda^\dagger}_{\text{Projection onto } \mathcal{R}(\Phi_\Lambda)}$$

$$P\Phi_\Lambda = 0$$

# Interference Cancellation

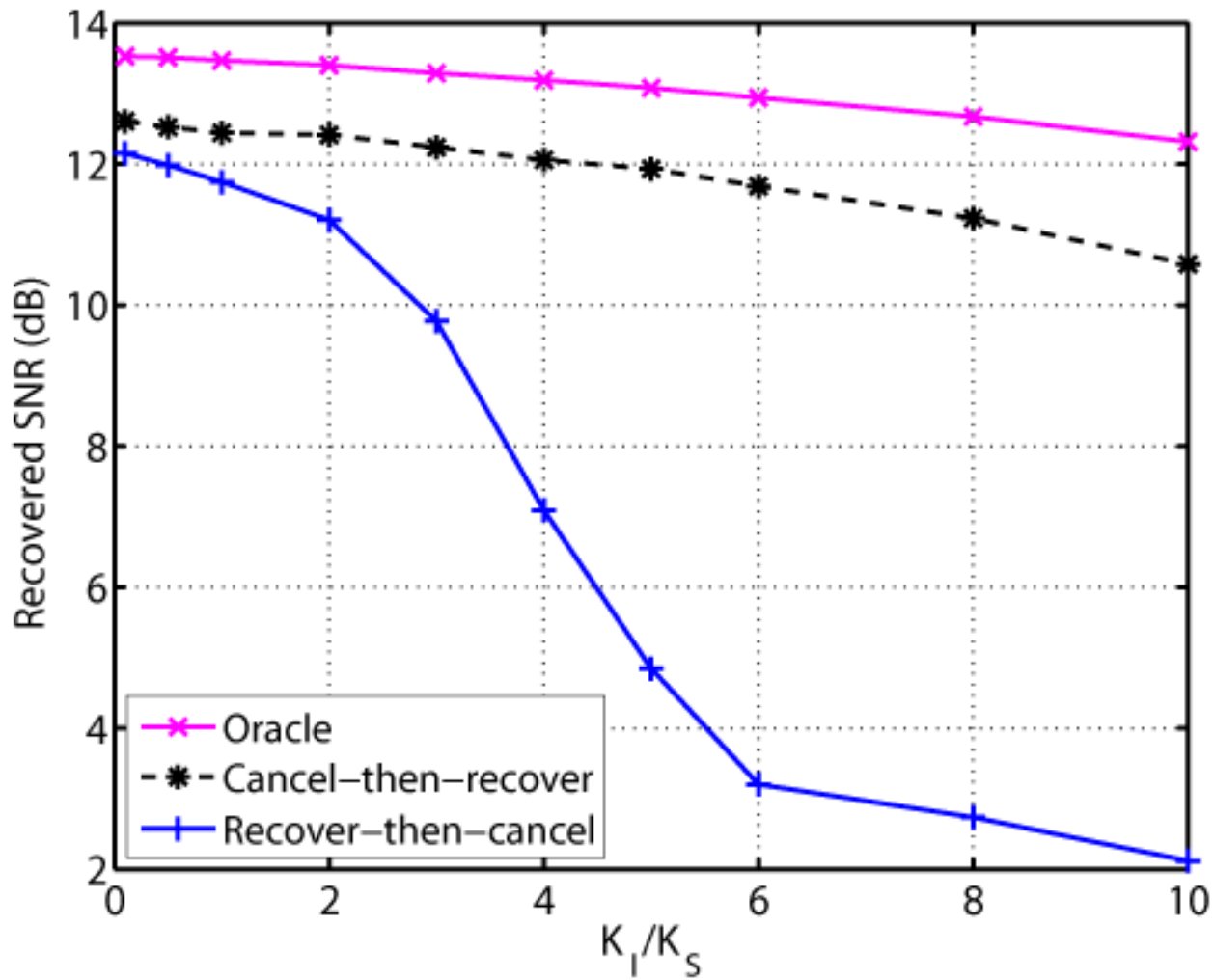
**Lemma:**

If  $\Phi$  satisfies the RIP of order  $S$ , then

$$\left(1 - \frac{\delta}{1 - \delta}\right) \|x\|_2^2 \leq \|P\Phi x\|_2^2 \leq (1 + \delta) \|x\|_2^2$$

provided that  $\|x\|_0 \leq S - |\Lambda|$  and  $\text{supp}(x) \cap \Lambda = \emptyset$ .

# Interference Cancellation in Action





# Interference Cancellation

## Lemma:

If  $\Phi$  satisfies the RIP of order  $S$ , then

$$\left(1 - \frac{\delta}{1 - \delta}\right) \|x\|_2^2 \leq \|P\Phi x\|_2^2 \leq (1 + \delta) \|x\|_2^2$$

provided that  $\|x\|_0 \leq S - |\Lambda|$  and  $\text{supp}(x) \cap \Lambda = \emptyset$ .

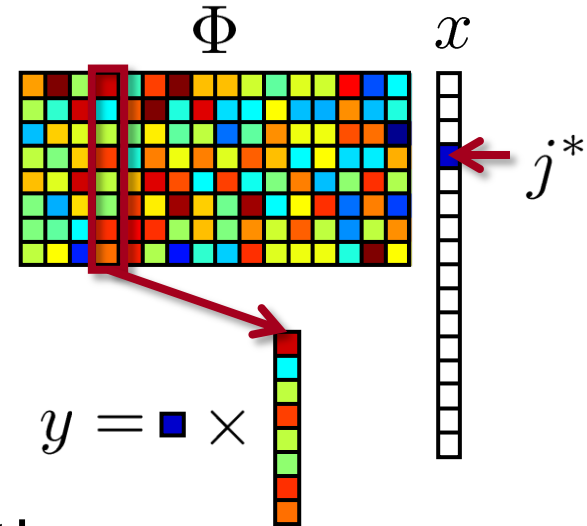
$$\longrightarrow |\langle Py, P\Phi_j \rangle - x_j| \leq \frac{\delta}{1 - \delta} \|x_{\Lambda^c}\|_2$$

# Aside: Orthogonal Matching Pursuit

OMP selects one index at a time

Iteration 1:

$$j^* = \arg \max_j |\langle y, \Phi_j \rangle|$$



If  $\Phi$  satisfies the RIP of order  $\|u \pm v\|_0$ , then

$$|\langle \Phi u, \Phi v \rangle - \langle u, v \rangle| \leq \delta \|u\|_2 \|v\|_2$$

Set  $u = x$  and  $v = e_j$

$$|\langle y, \Phi_j \rangle - x_j| \leq \delta \|x\|_2$$

# Aside: Orthogonal Matching Pursuit

Subsequent Iterations:

$$j^* = \arg \max_j |\langle Py, P\Phi_j \rangle|$$

$$P = I - \Phi_\Lambda \Phi_\Lambda^\dagger$$

$$P\Phi_\Lambda = 0 \quad \longrightarrow \quad P\Phi x = P\Phi x_{\Lambda^c}$$

$$\longrightarrow \quad |\langle Py, P\Phi_j \rangle - x_j| \leq \frac{\delta}{1 - \delta} \|x_{\Lambda^c}\|_2$$

# Aside: Orthogonal Matching Pursuit

## Theorem:

Suppose  $x$  is  $S$ -sparse and  $y = \Phi x$ .

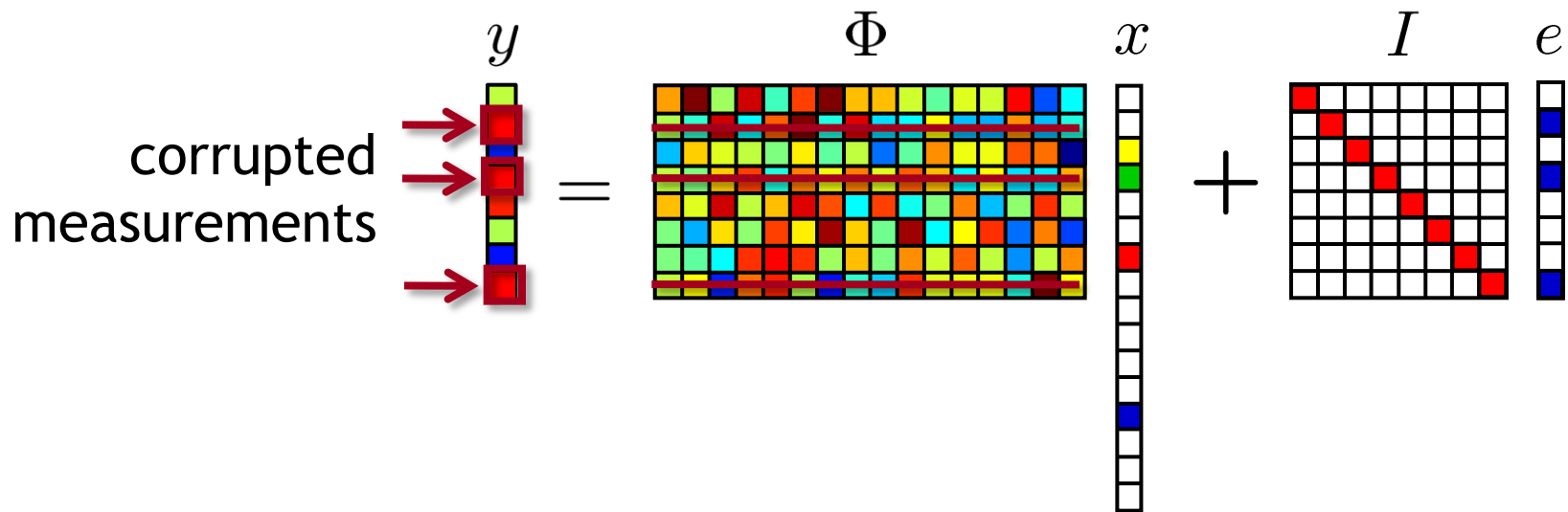
If  $\Phi$  satisfies the RIP of order  $S + 1$  with constant  $\delta < \frac{1}{3\sqrt{S}}$ , then the  $j^*$  identified at each iteration will be a nonzero entry of  $x$ .

➡ Exact recovery after  $S$  iterations.

Argument provides simplified proofs for other orthogonal greedy algorithms (e.g. ROMP) that are robust to noise

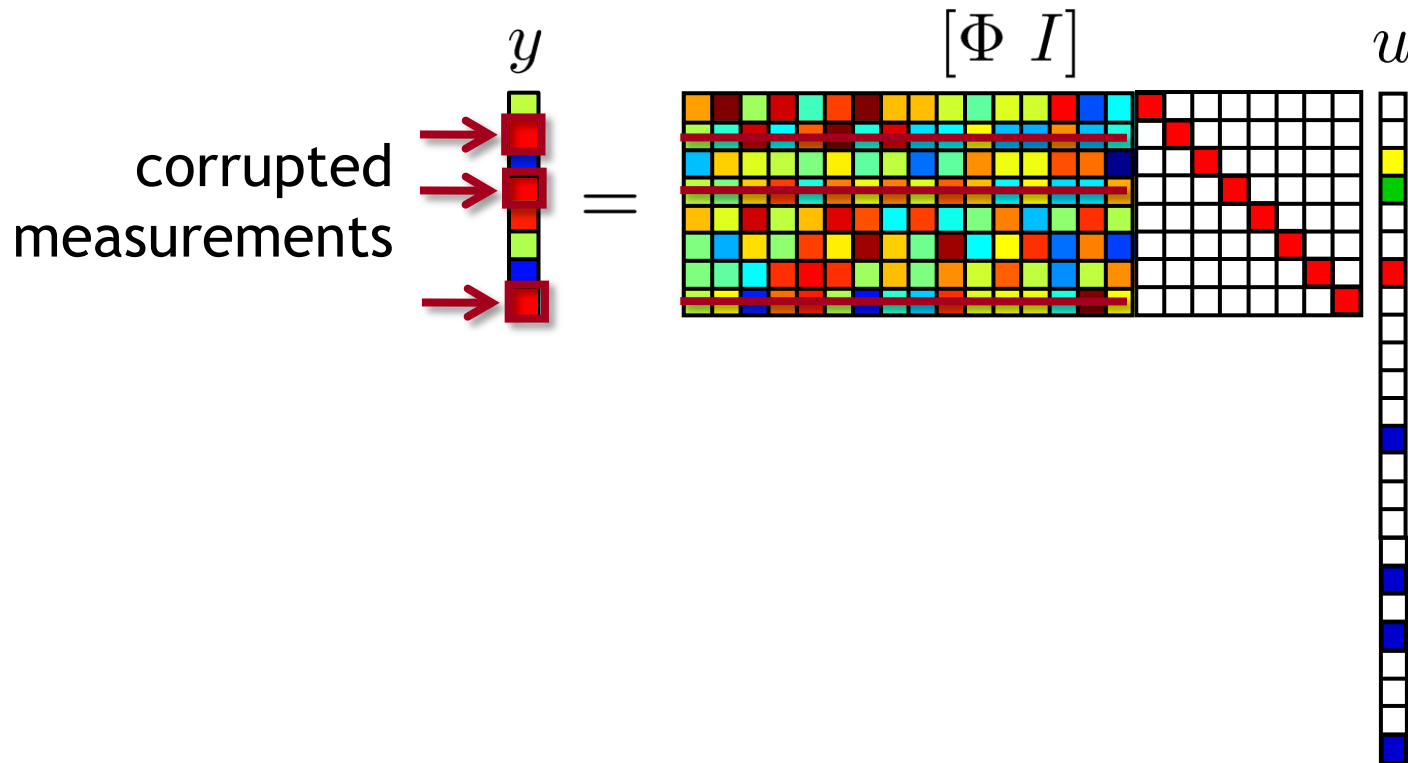
# Measurement Interference Cancellation

What about structured measurement noise?



# Measurement Interference Cancellation

What about structured measurement noise?



# Justice Pursuit

$$\begin{aligned} \hat{u} &= \arg \min_u \|u\|_1 \\ \text{s.t. } & y = [\Phi \ I] u \end{aligned}$$

Does this matrix satisfy the RIP?

$$\|[\Phi \ I]u\|_2^2 = \|\Phi x\|_2^2 + 2e^T \Phi x + \|e\|_2^2 \approx \|x\|_2^2 + \|e\|_2^2$$

## Theorem:

If  $\Phi$  is a sub-Gaussian matrix with

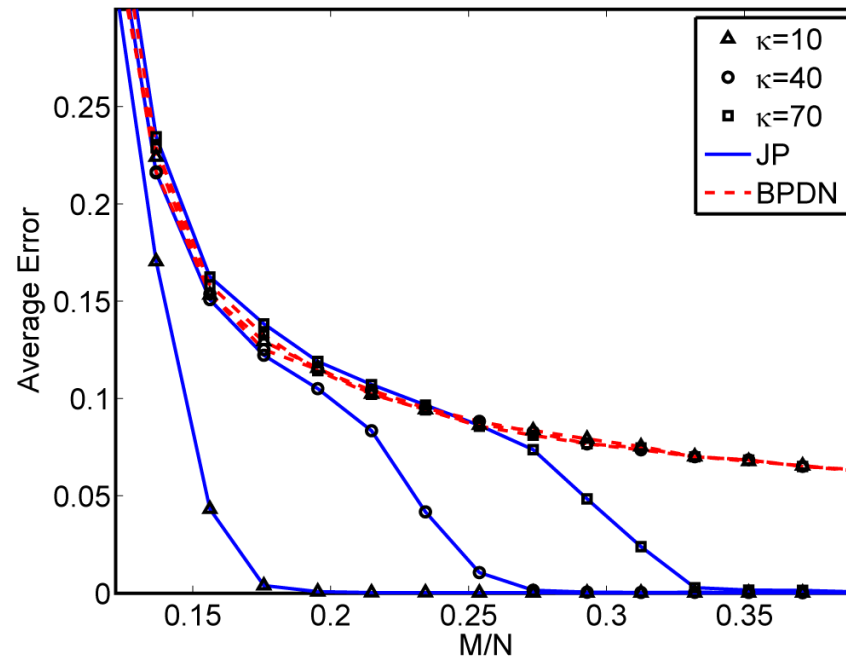
$$M = O \left( (S + \kappa) \log \left( \frac{N + M}{S + \kappa} \right) \right)$$

then  $[\Phi \ I]$  satisfies the RIP of order  $(S + \kappa)$  with probability at least  $1 - 3e^{-CM}$ .

# Justice Pursuit

We can recover sparse signals *exactly* in the presence of *unbounded* sparse noise

Fixed  $\|e\|_2 = 0.1$



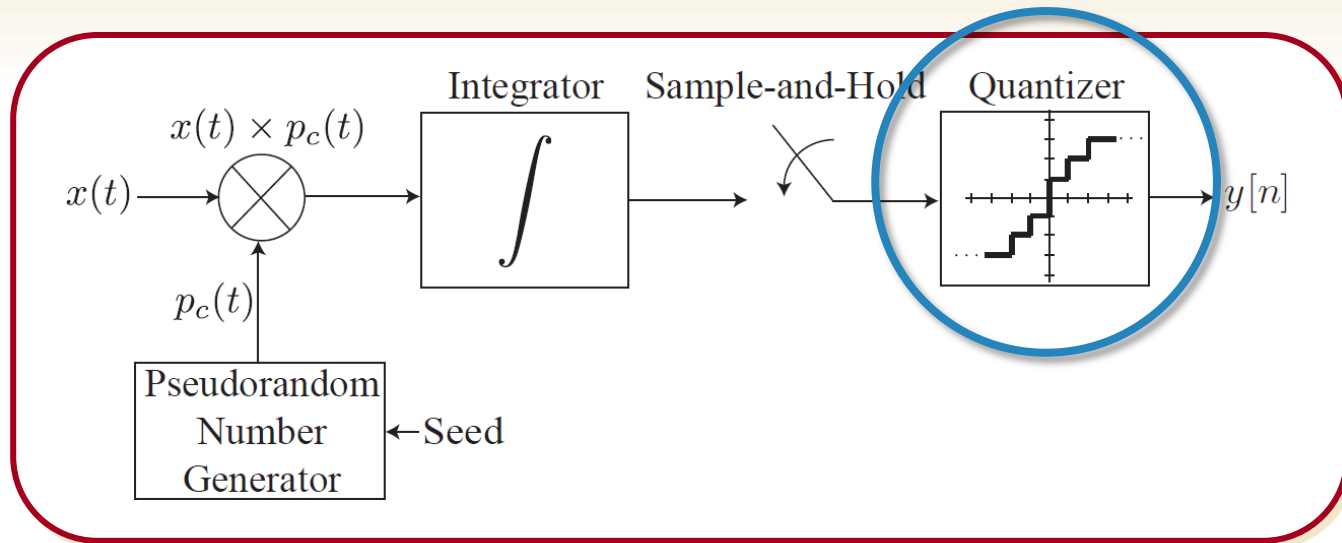


# Conclusions

- CS systems are sensitive to noisy signals
  - if our input signal is very noisy, it isn't really very sparse
  - when noise is large, *measurements matter*
  - exploit sparsity in a different manner - e.g., adaptivity
- CS can be highly robust to *interference*
  - structured signal noise
  - structured measurement noise
- What about quantization noise?

# Quantization Noise

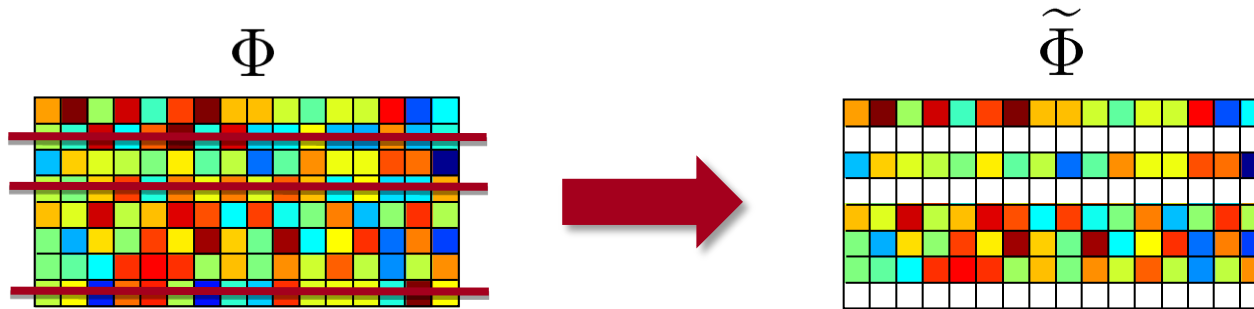
# Signal Recovery with Quantization



- Finite-range quantization leads to *saturation* and *unbounded errors*
- Quantization noise changes as we change the sampling rate

# Saturation Strategies

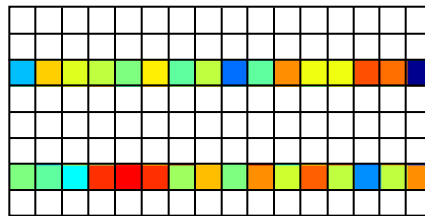
- **Rejection:** Ignore saturated measurements



- **Consistency:** Retain saturated measurements. Use them only as inequality constraints on the recovered signal
- If the rejection approach works, the consistency approach should automatically do better

# Rejection and Democracy

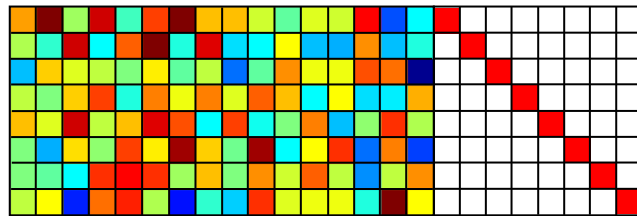
- The RIP is *not sufficient* for the rejection approach
- Example:  $\Phi = I$ 
  - perfect isometry
  - *every* measurement must be kept
- We would like to be able to say that *any* submatrix of  $\Phi$  with sufficiently many rows will still satisfy the RIP



- Strong, *adversarial* form of “democracy”

# Sketch of Proof

- Step 1: Concatenate the identity to  $\Phi$



## Theorem:

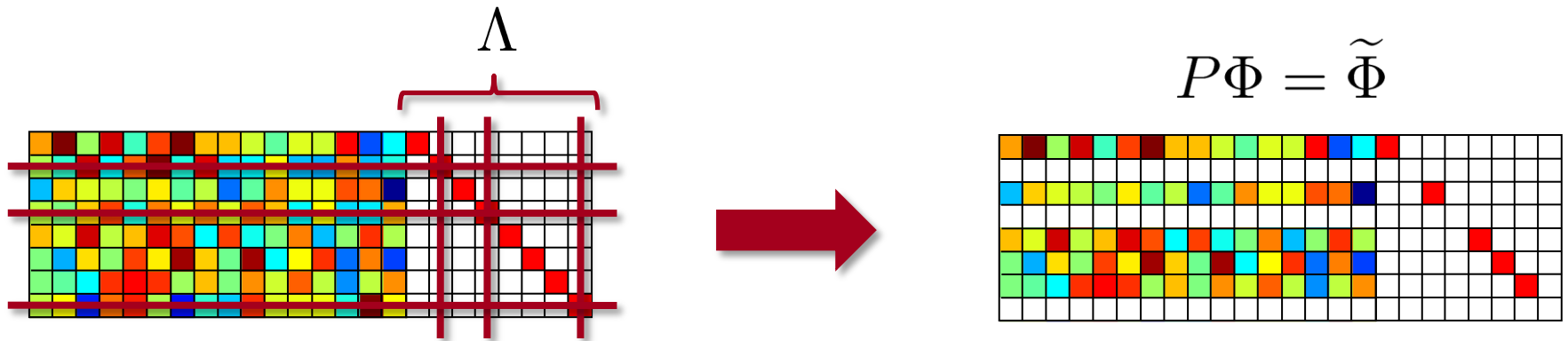
If  $\Phi$  is a sub-Gaussian matrix with

$$M = O \left( S \log \left( \frac{N + M}{S} \right) \right)$$

then  $[\Phi \ I]$  satisfies the RIP of order  $S$  with probability at least  $1 - 3e^{-CM}$ .

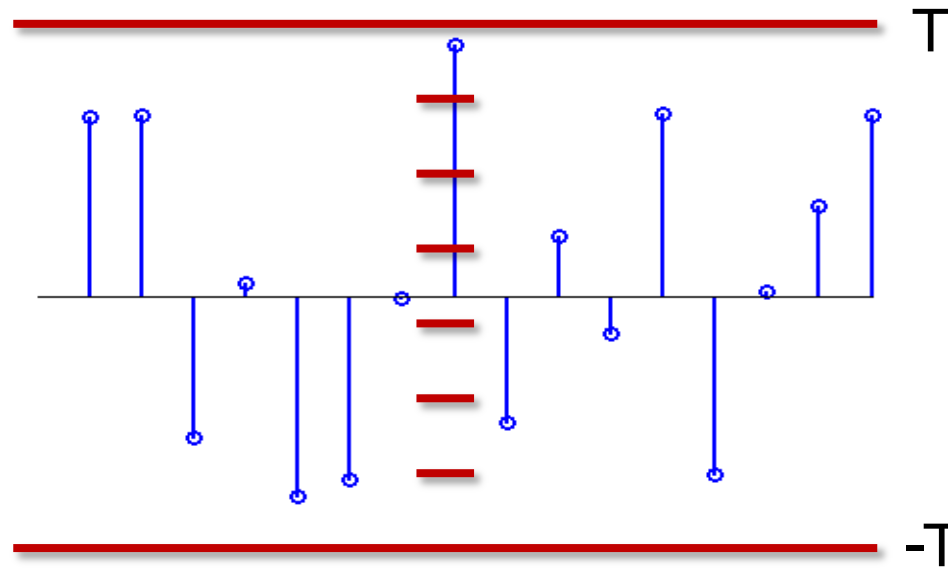
# Sketch of Proof

- Step 2: Combine with the “interference cancellation” lemma



- The fact that  $[\Phi \ I]$  satisfies the RIP implies that if we take  $D$  extra measurements, then we can delete  $O(D)$  arbitrary rows of  $\Phi$  and retain the RIP
- This is a strong *adversarial* notion of democracy

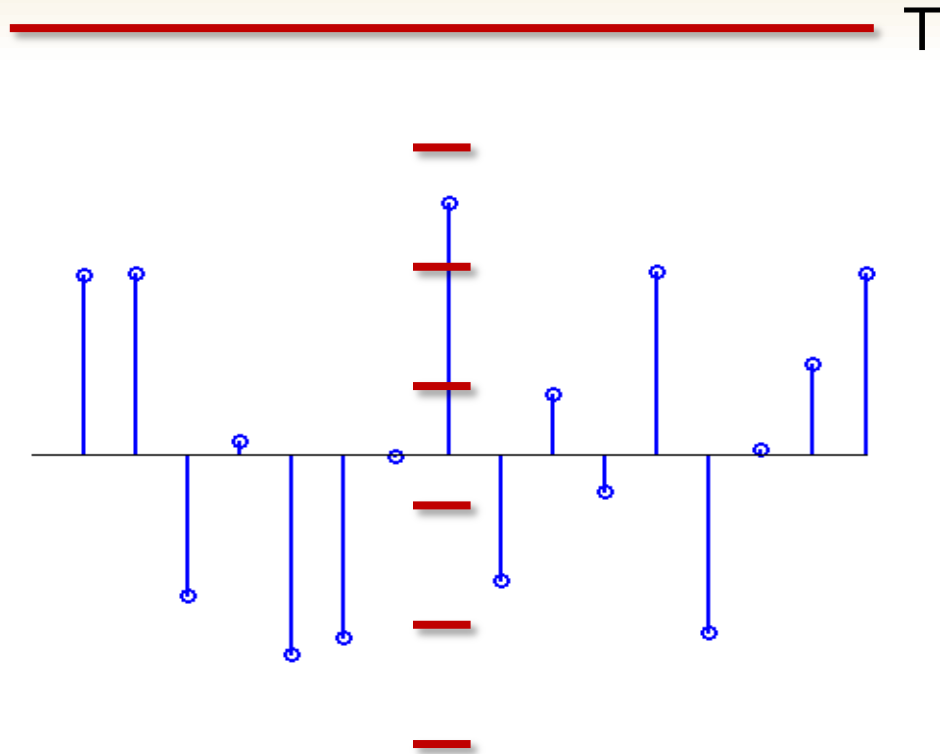
# Rejection In Practice



$$\text{SNR} = 10 \log_{10} \left( \frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right)$$

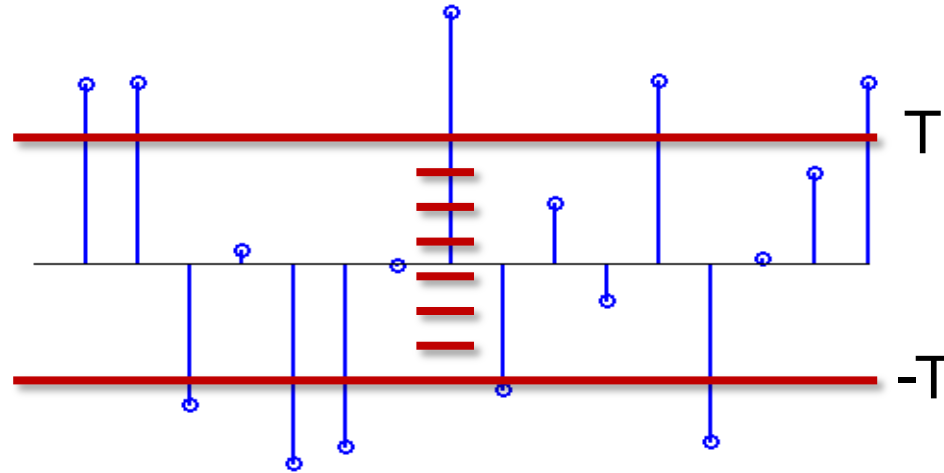


# Rejection In Practice



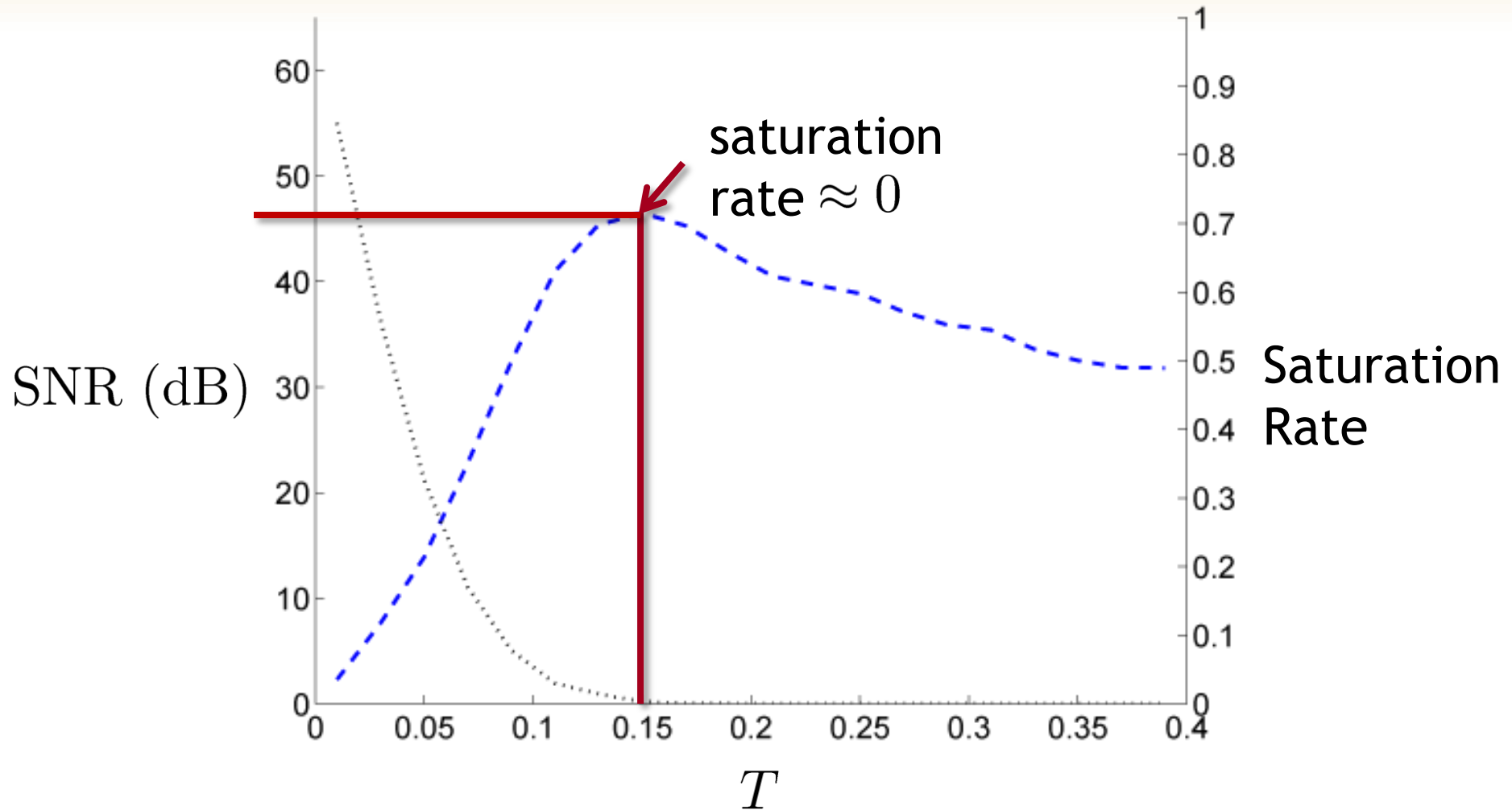
$$\text{SNR} = 10 \log_{10} \left( \frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right)$$

# Rejection In Practice

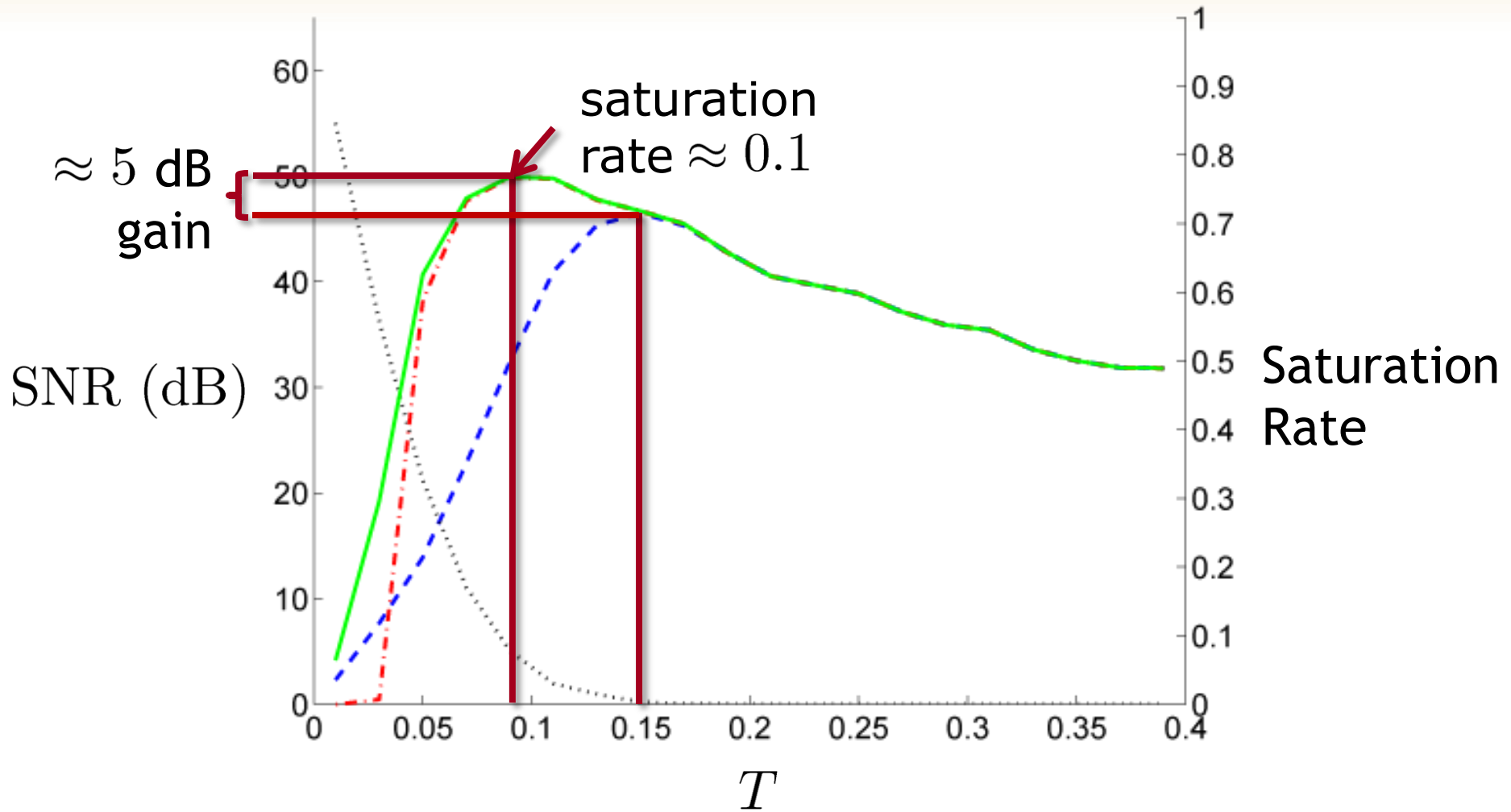


$$\text{SNR} = 10 \log_{10} \left( \frac{\|x\|_2^2}{\|\hat{x} - x\|_2^2} \right)$$

# Benefits of Saturation

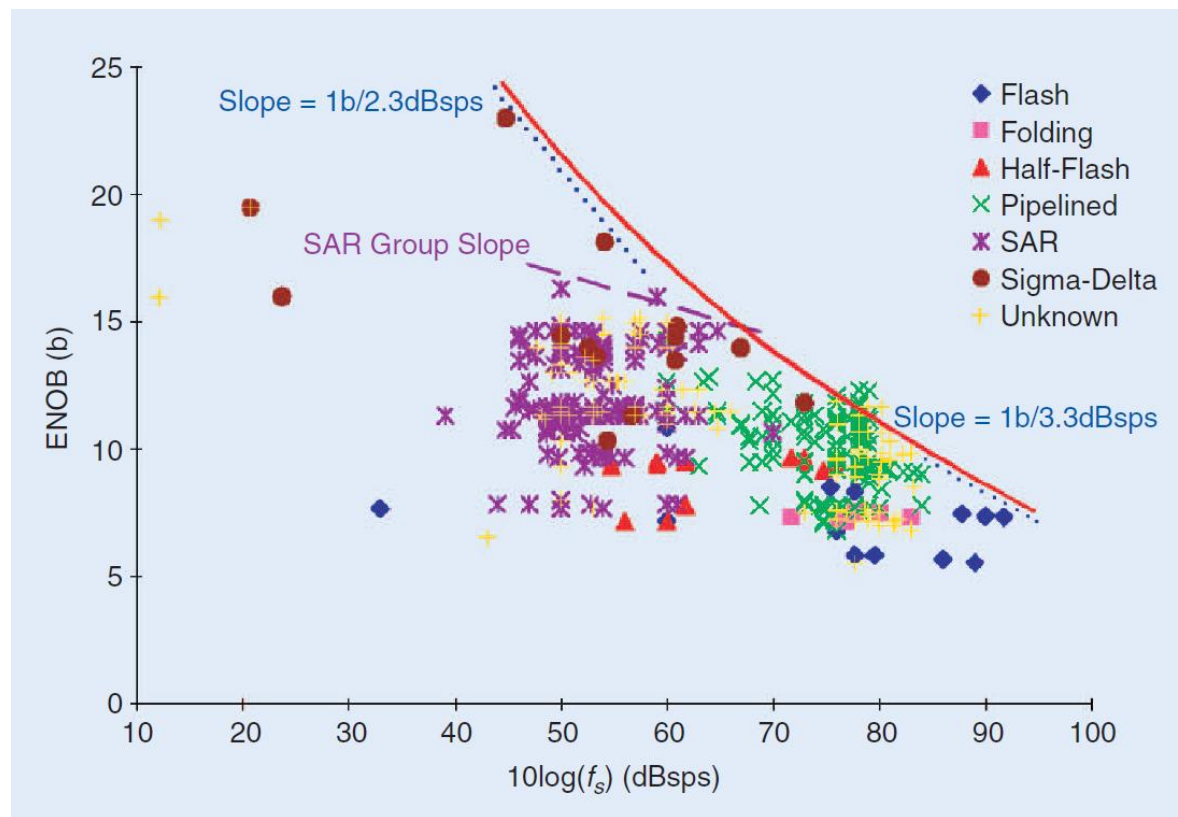


# Benefits of Saturation

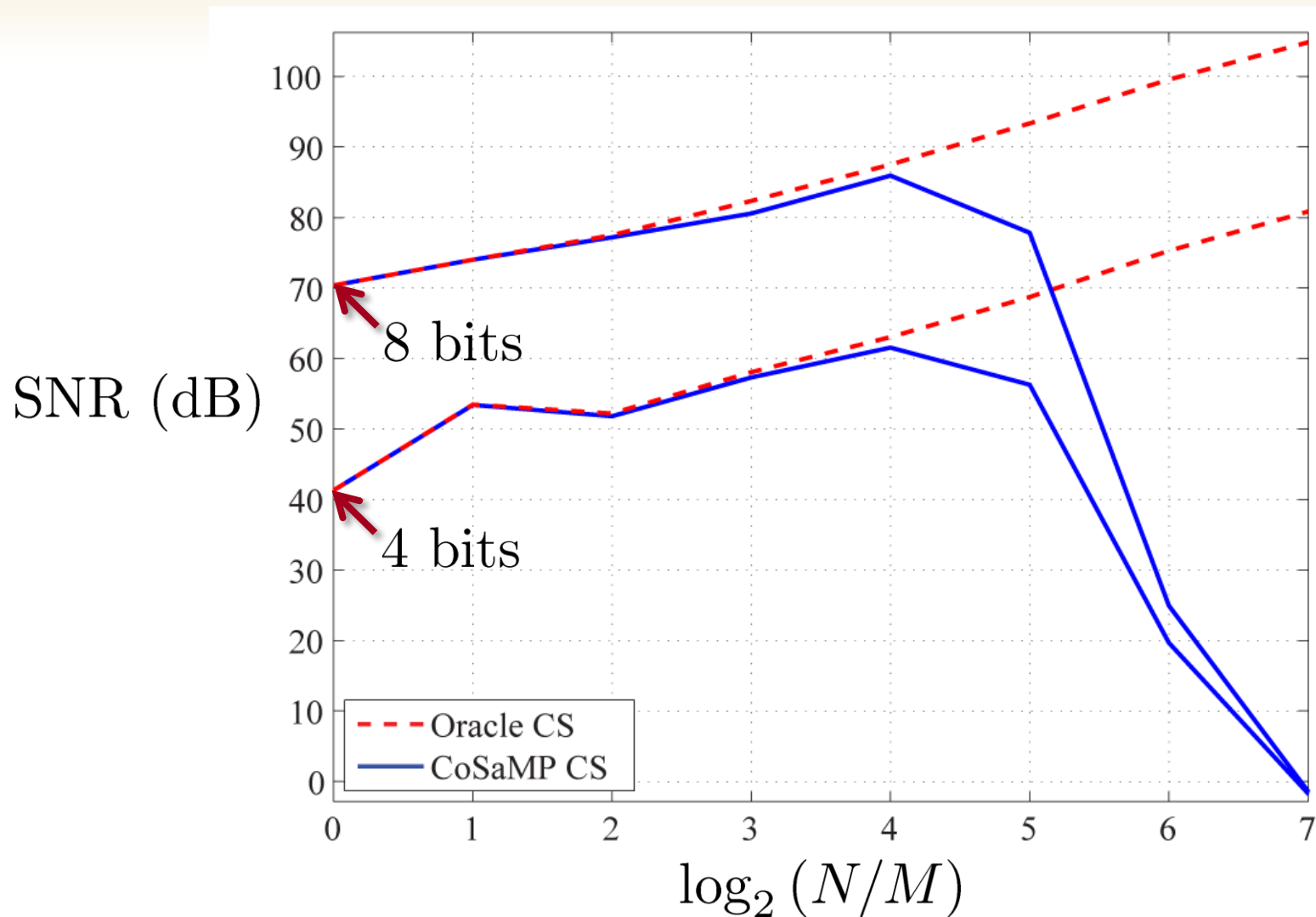


# Potential for SNR Improvement?

By sampling at a lower rate, we can quantize to a higher bit-depth, allowing for potential gains



# Empirical SNR Improvement



# Conclusions

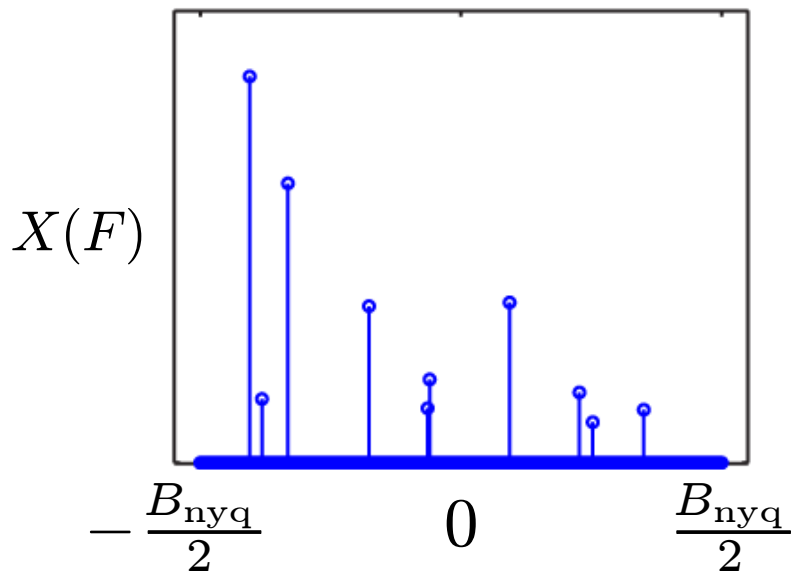
- CS is robust to quantization noise in a non-traditional sense
- Democracy is a major advantage of CS measurements
- CS offers the potential to significantly boost dynamic range
  - can offset drawbacks associated with noise
- When is CS most useful?
  - performance is limited by quantization (high bandwidth apps)
  - when your signal is sparse (not too noisy)

# **Real-World Signal Models**



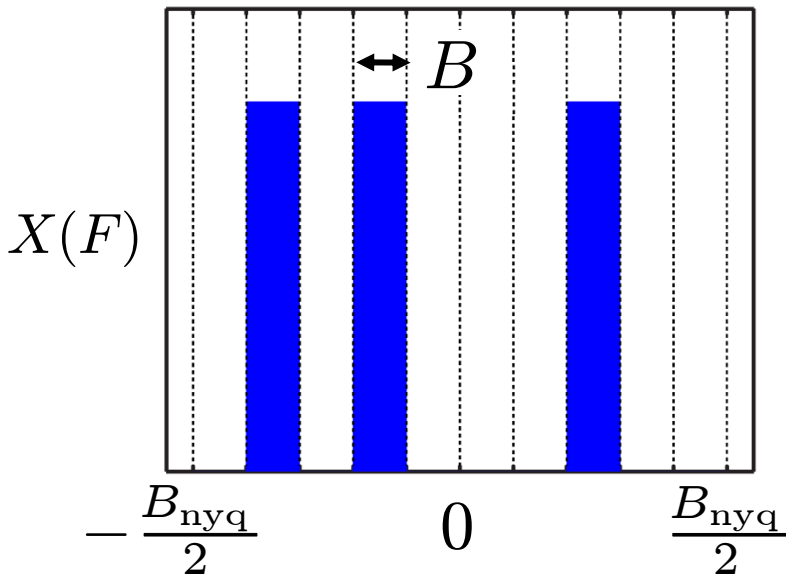
# Candidate Analog Signal Models

	Model for $x(t)$	Basis for $x$	Sparsity level for $x$
multitone	sum of $S$ “on-grid” tones	$\Psi = \text{DFT}$	$S$ -sparse



# Candidate Analog Signal Models

	Model for $x(t)$	Basis for $x$	Sparsity level for $x$
multitone	sum of $S$ “on-grid” tones	$\Psi = \text{DFT}$	$S$ -sparse
multiband	$K$ occupied bands of bandwidth $B$	$\Psi = ?$	?



- Landau
- Bresler, Feng, Venkataramani
- Eldar, Mishali

# The Problem with the DFT

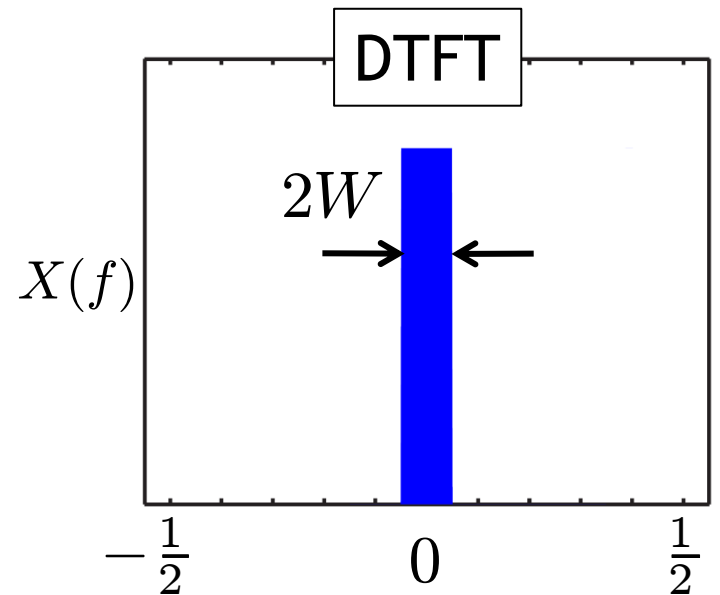
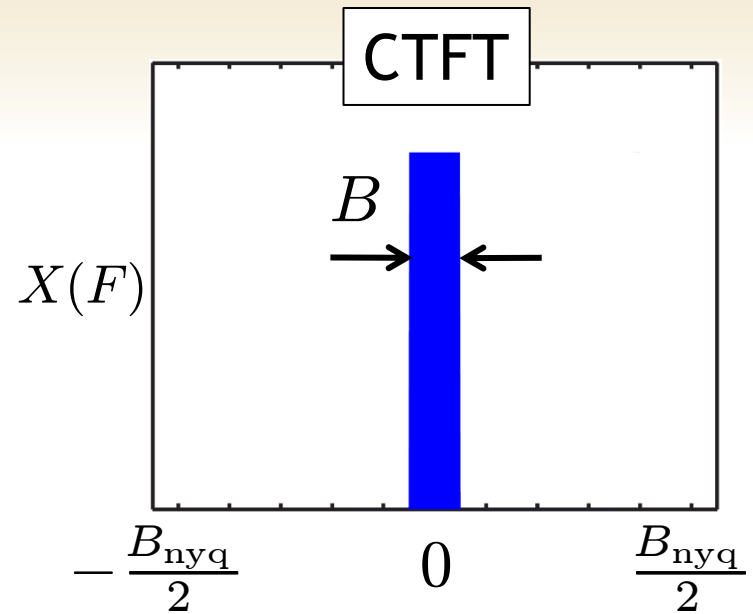
$$x(t) = \int_{-\frac{B}{2}}^{\frac{B}{2}} X(F) e^{j2\pi Ft} dF$$



sampling

$$x[n] = \int_{-W}^W X(f) e^{j2\pi fn} df, \quad \forall n$$

$$W = \frac{B}{2B_{\text{nyq}}}$$



# The Problem with the DFT

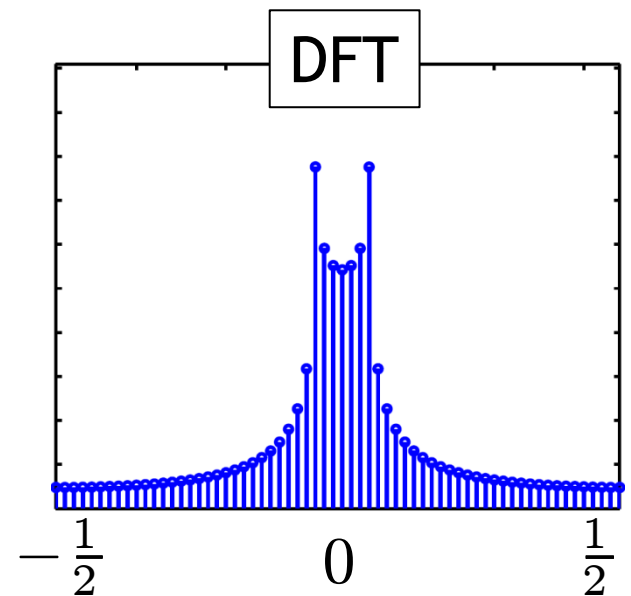
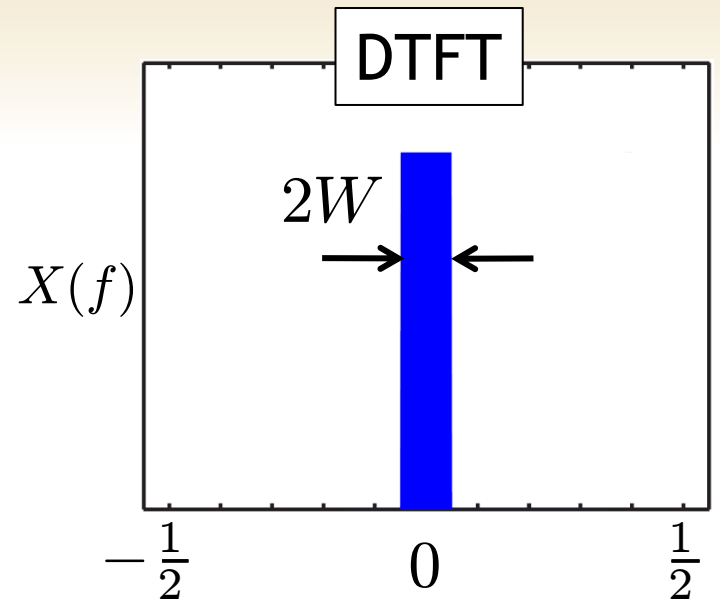
$$x[n] = \int_{-W}^W X(f) e^{j2\pi f n} df, \quad \forall n$$



time-limiting

$$x = \sum_{k=0}^{N-1} X_k e^{\frac{k}{N}}, \quad e_f := \begin{bmatrix} e^{j2\pi f 0} \\ e^{j2\pi f} \\ \vdots \\ e^{j2\pi f(N-1)} \end{bmatrix}$$

**NOT SPARSE**

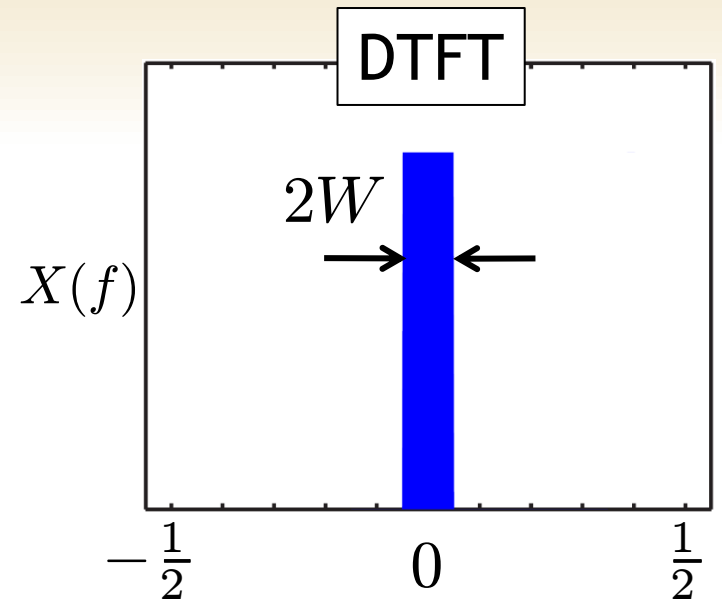


# Alternative Perspective

$$x[n] = \int_{-W}^W X(f) e^{j2\pi f n} df, \quad \forall n$$



time-limiting



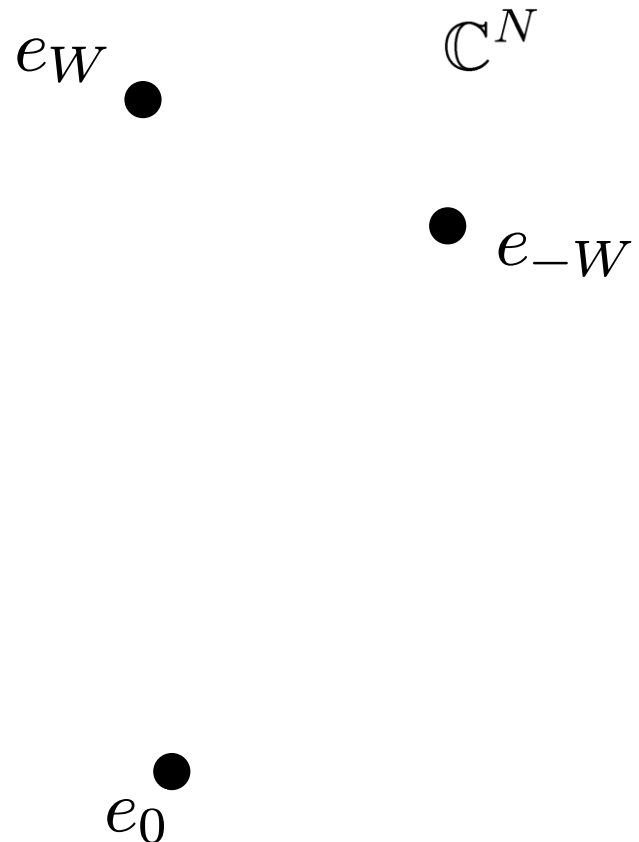
$$\mathcal{T}_N(x[n]) = \int_{-W}^W X(f) \mathcal{T}_N(e^{j2\pi f n}) df, \quad \forall n$$

# Building Blocks for Lowpass Signals

Time-limited complex exponentials form a “basis” for bandlimited signals

$$x = \int_{-W}^W X(f) e_f df$$

$$e_f := \begin{bmatrix} e^{j2\pi f0} \\ e^{j2\pi f} \\ \vdots \\ e^{j2\pi f(N-1)} \end{bmatrix}$$



The problem: we need infinitely many of them.

# Best Subspace Fit

Suppose that we wish to minimize

$$\int_{-W}^W \|e_f - P_Q e_f\|_2^2 df$$

over all subspaces  $Q$  of dimension  $k$ .

Optimal subspace is spanned  
by the first  $k$  “DPSS vectors”.

# Discrete Prolate Spheroidal Sequences (DPSS's)

*Slepian [1978]*: Given an integer  $N$  and  $W \leq \frac{1}{2}$ , the DPSS's are a collection of  $N$  vectors

$$s_0, s_1, \dots, s_{N-1} \in \mathbb{R}^N$$

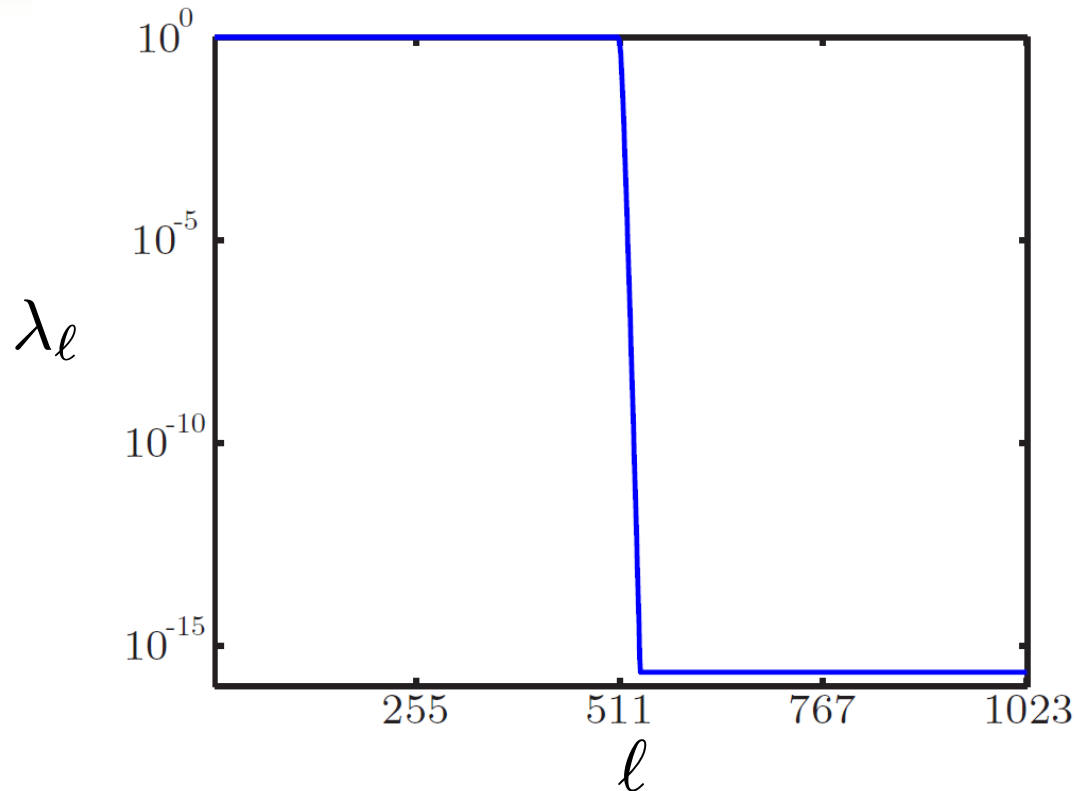
that satisfy

$$\mathcal{T}_N(\mathcal{B}_W(s_\ell)) = \lambda_\ell s_\ell.$$

The DPSS's are perfectly time-limited, but when  $\lambda_\ell \approx 1$  they are highly concentrated in frequency.



# DPSS Eigenvalue Concentration



$$N = 1024$$

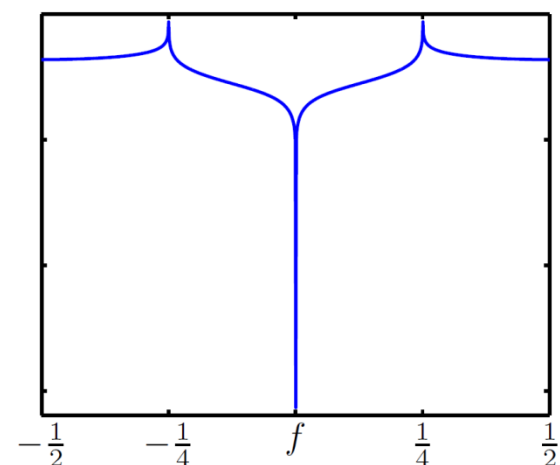
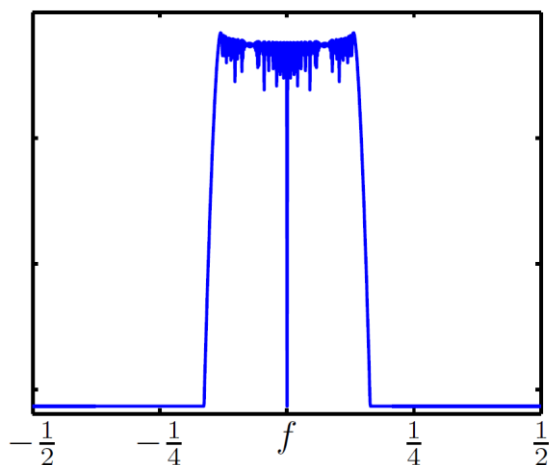
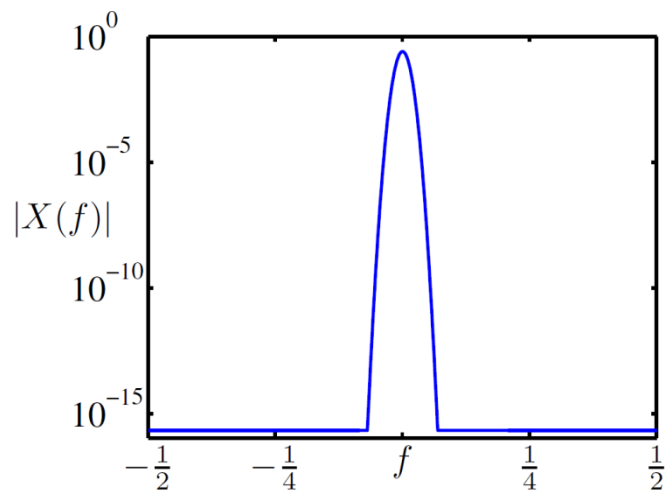
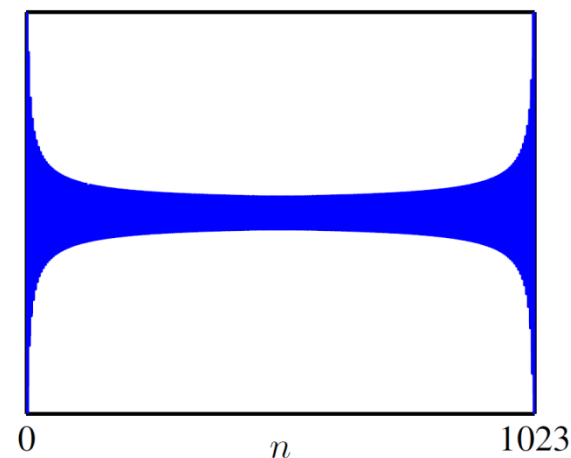
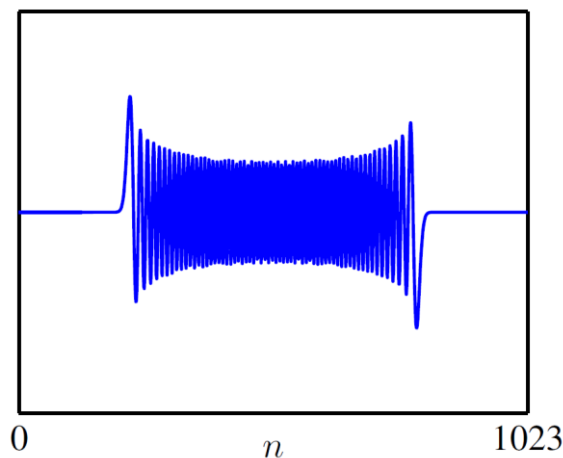
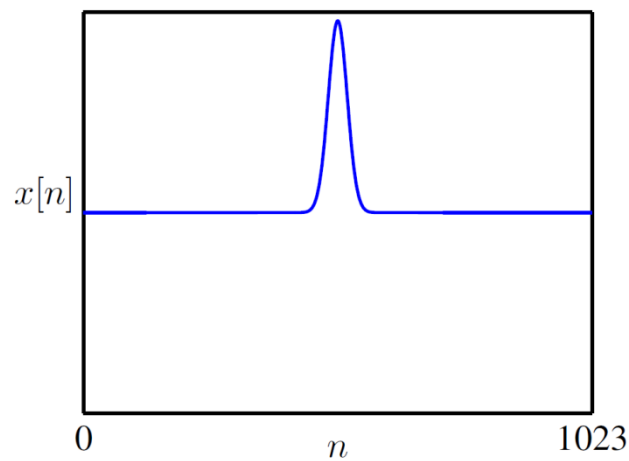
$$W = \frac{1}{4}$$

$$2NW = 512$$

The first  $\approx 2NW$  eigenvalues  $\approx 1$ .  
The remaining eigenvalues  $\approx 0$ .

# DPSS Examples

$$N = 1024 \quad W = \frac{1}{4}$$



$$\ell = 0$$

$$\ell = 127$$

$$\ell = 511$$

# Recall: Best Subspace Fit

Suppose that we wish to minimize

$$\int_{-W}^W \|e_f - P_Q e_f\|_2^2 df$$

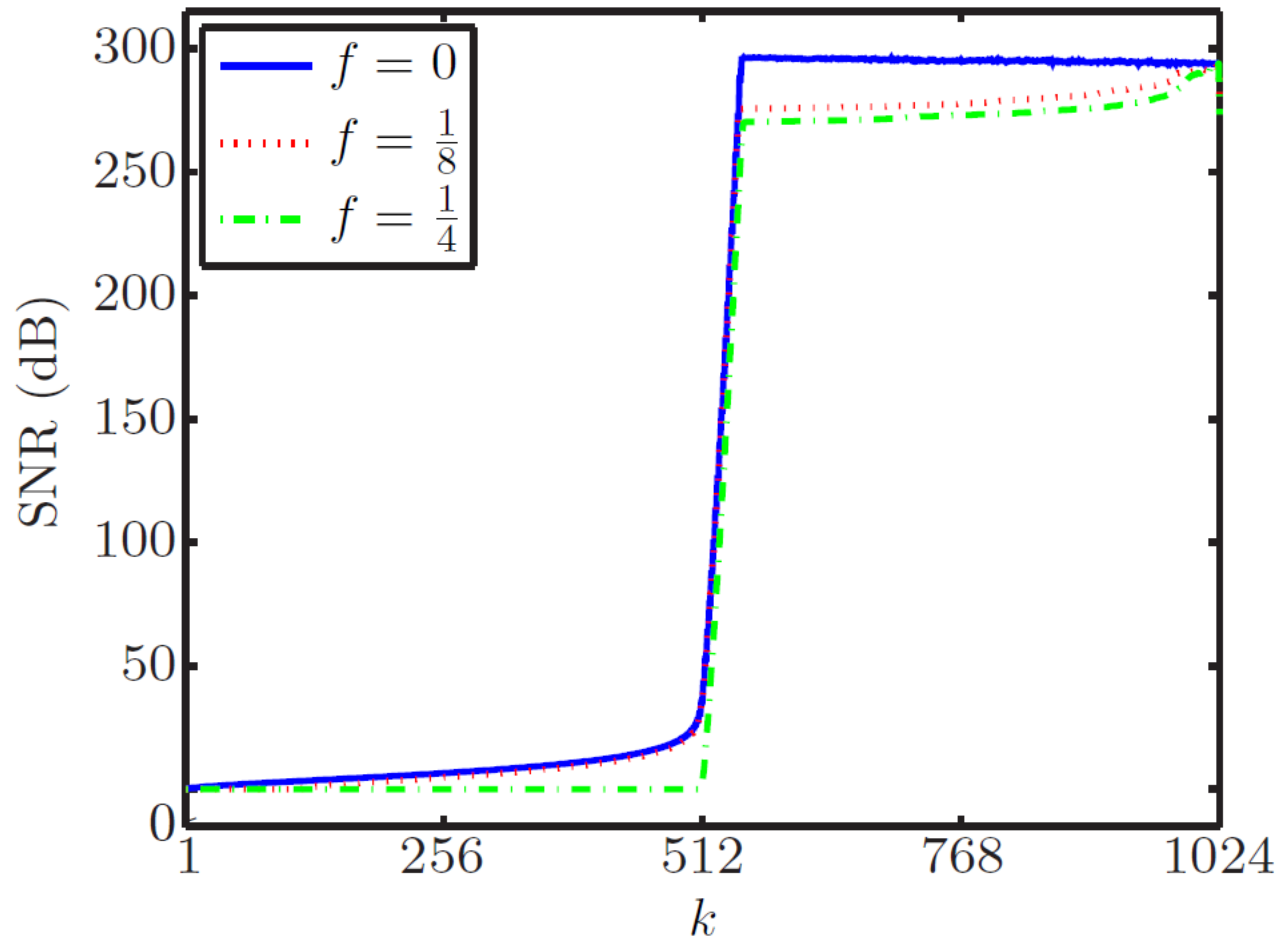
over all subspaces  $Q$  of dimension  $k$ .

Optimal subspace is spanned by the first  $k$  “DPSS vectors”.

$$\int_{-W}^W \|e_f - P_Q e_f\|_2^2 df = \sum_{\ell=k}^{N-1} \lambda_\ell$$

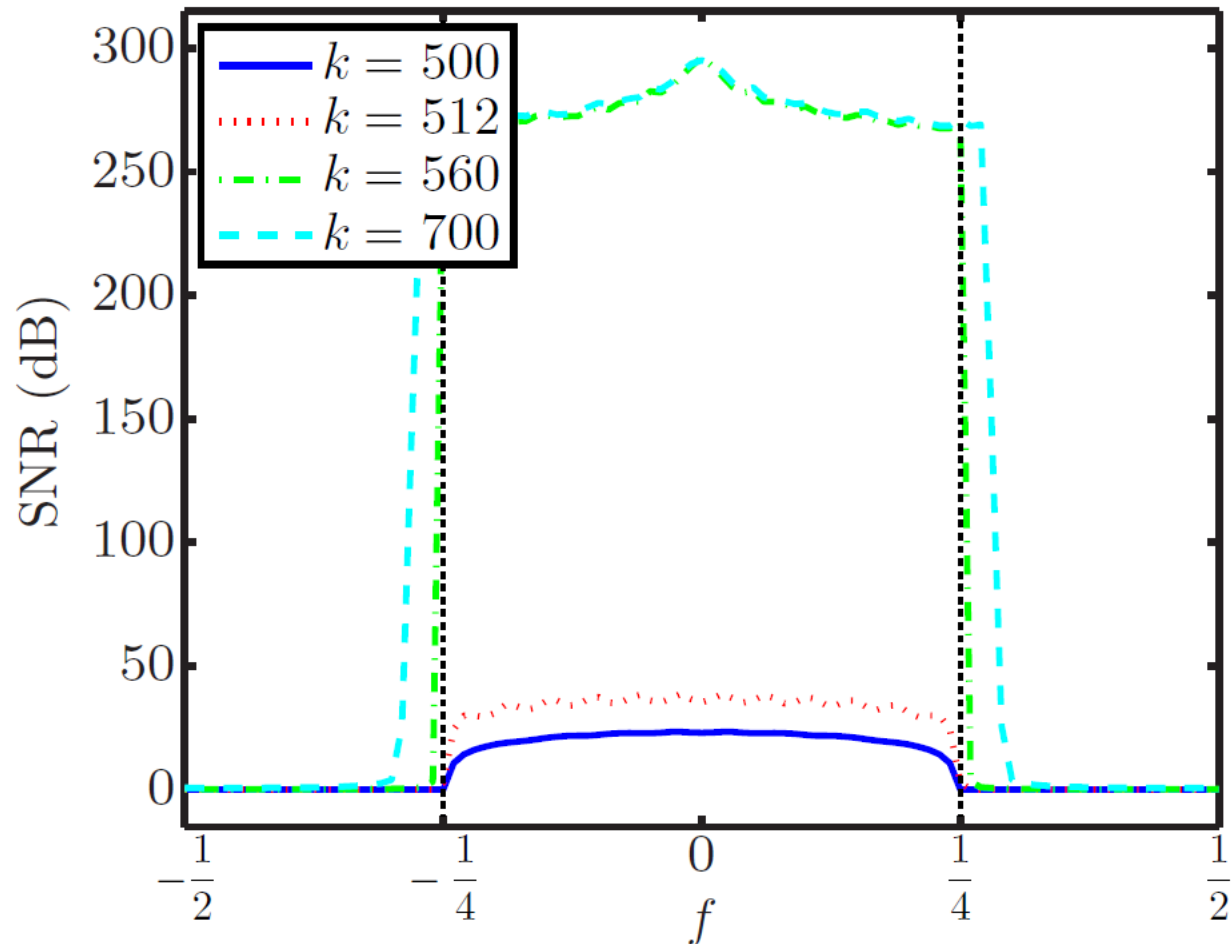
# Approximation of Bandlimited Signals

$$\text{SNR} = 20 \log_{10} \left( \frac{\|e_f\|}{\|e_f - P_Q e_f\|} \right) \text{ dB}$$

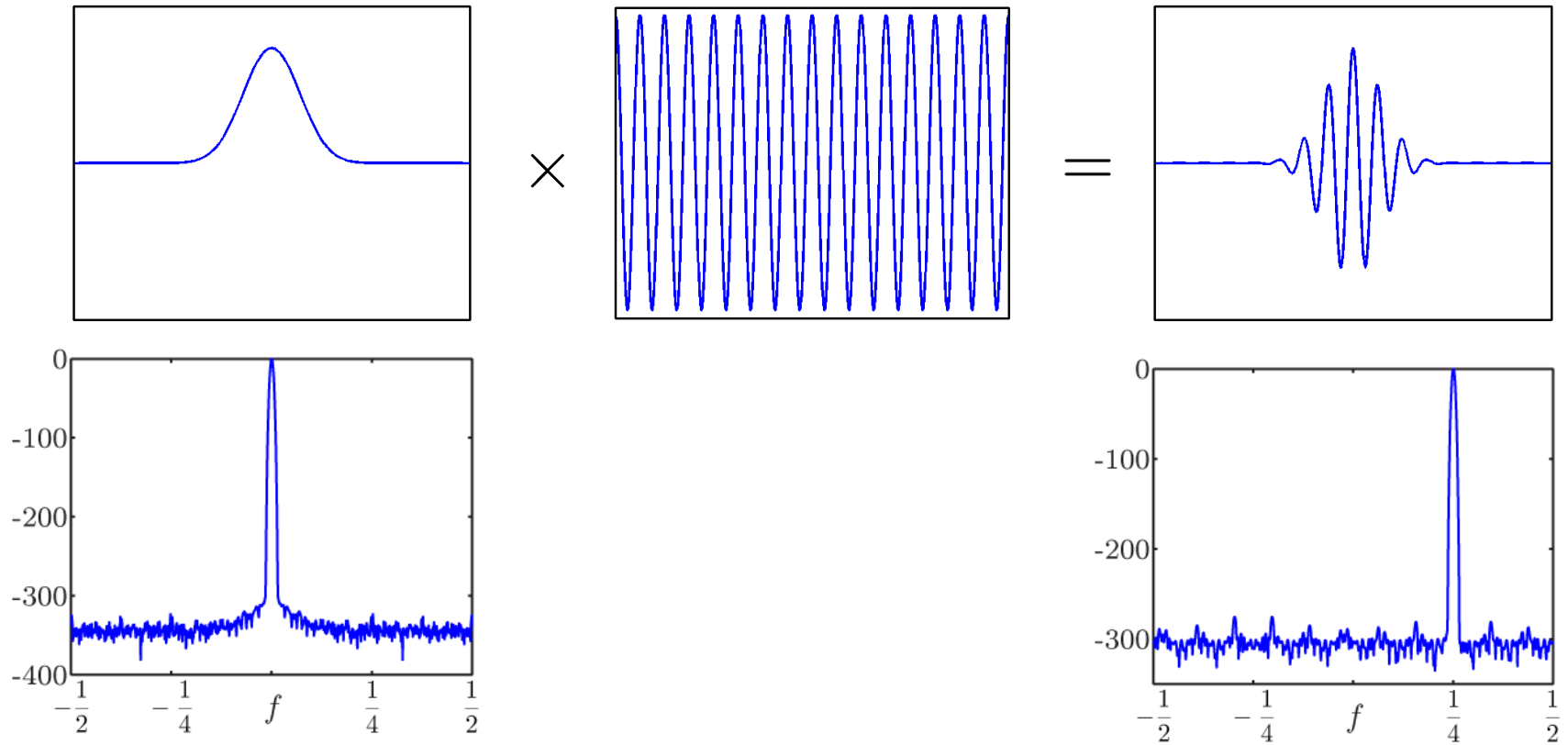


# Approximation of Bandlimited Signals

$$\text{SNR} = 20 \log_{10} \left( \frac{\|e_f\|}{\|e_f - P_Q e_f\|} \right) \text{ dB}$$



# DPSS's for Bandpass Signals

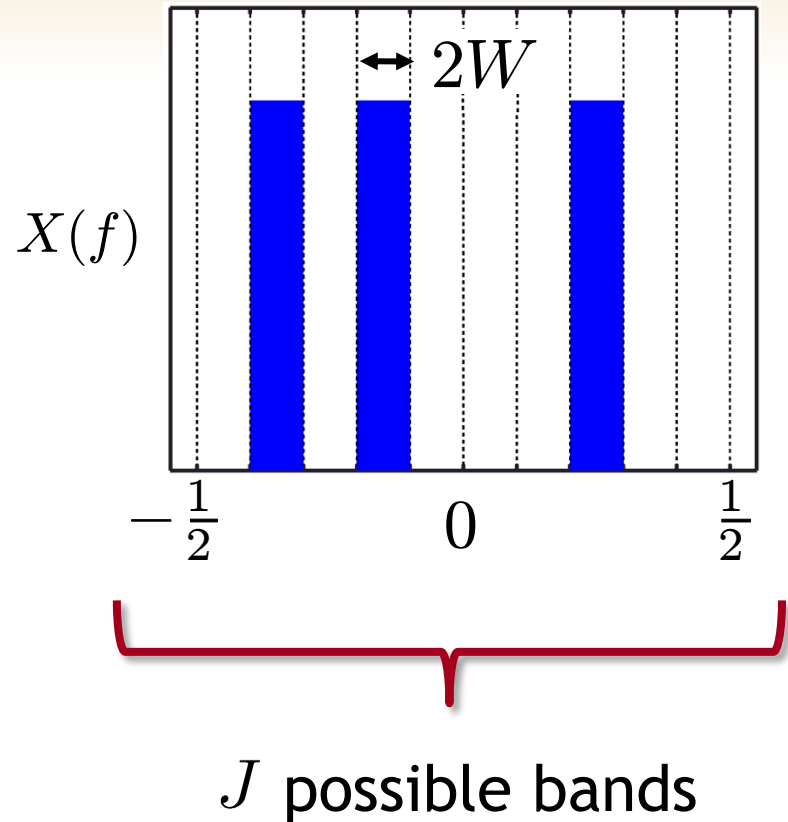


# DPSS Dictionaries for CS

Modulate  $k$  DPSS vectors  
to center of each band:

$$\Psi = [\Psi_1, \Psi_2, \dots, \Psi_J]$$

approximately square  
if  $k \approx 2NW$



Most multiband signals, when sampled and time-limited,  
are well-approximated by a sparse representation in  $\Psi$ .

# DPSS Dictionaries and the RIP

## Theorem:

Suppose that  $\Phi$  is sub-Gaussian and that the  $\Psi_i$  are constructed with  $k = (1 - \epsilon)2NW$ . If

$$M \geq CS \log(N/S)$$

then with high probability  $\Phi\Psi$  will satisfy the RIP of order  $S$ .

$K$  occupied bands  $\rightarrow S \approx KNB/B_{\text{nyq}}$

$$\frac{M}{N} \geq C' \frac{KB}{B_{\text{nyq}}} \log \left( \frac{B_{\text{nyq}}}{KB} \right)$$



# Block-Sparse Recovery

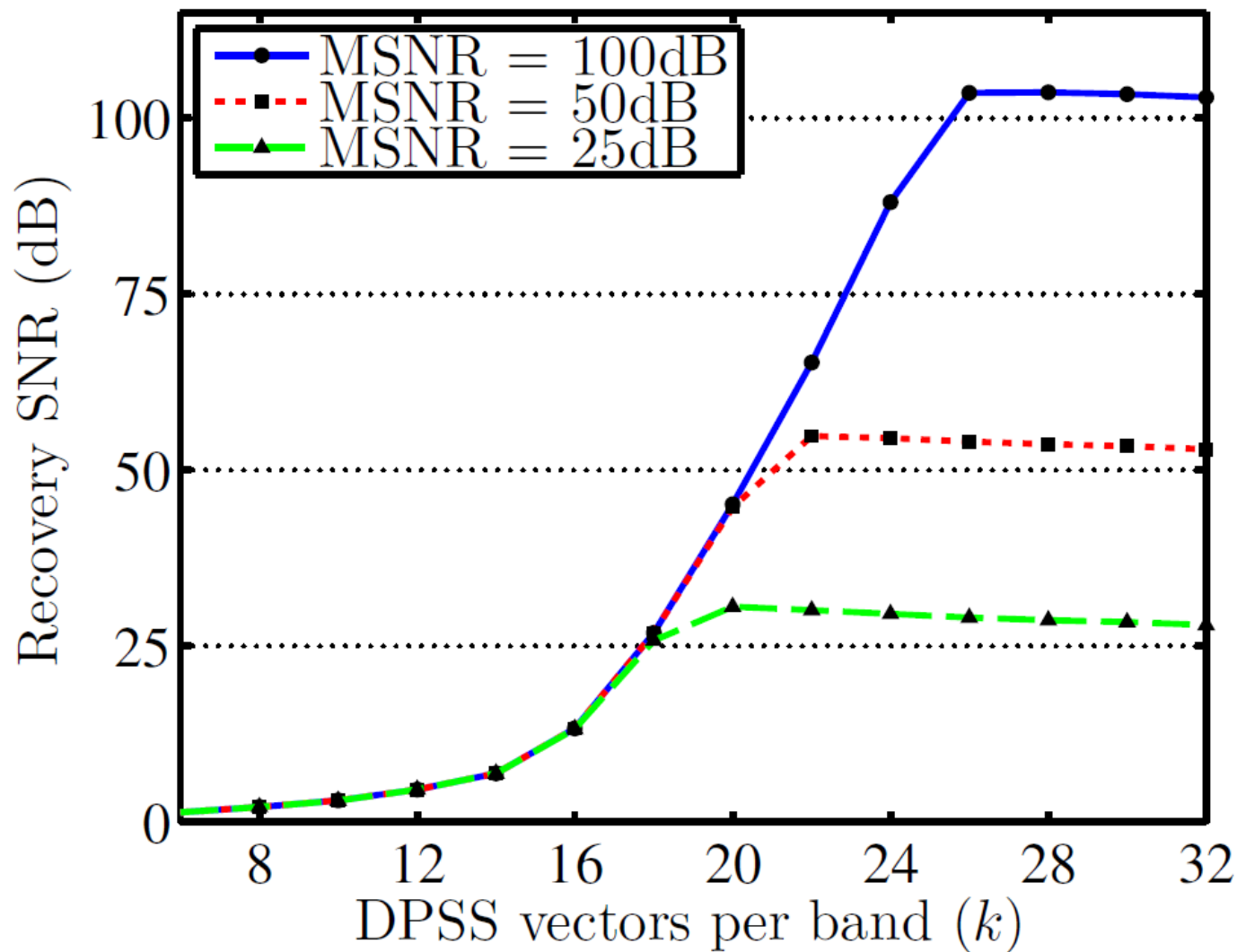
Nonzero coefficients of  $\alpha$  should be clustered in blocks according to the occupied frequency bands

$$x = [\Psi_1, \Psi_2, \dots, \Psi_J] \begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_J \end{bmatrix}$$

This can be leveraged to reduce the required number of measurements and improve performance through “model-based CS”

- Baraniuk et al. [2008, 2009, 2010]
- Blumensath and Davies [2009, 2011]

# Empirical Results: Noise



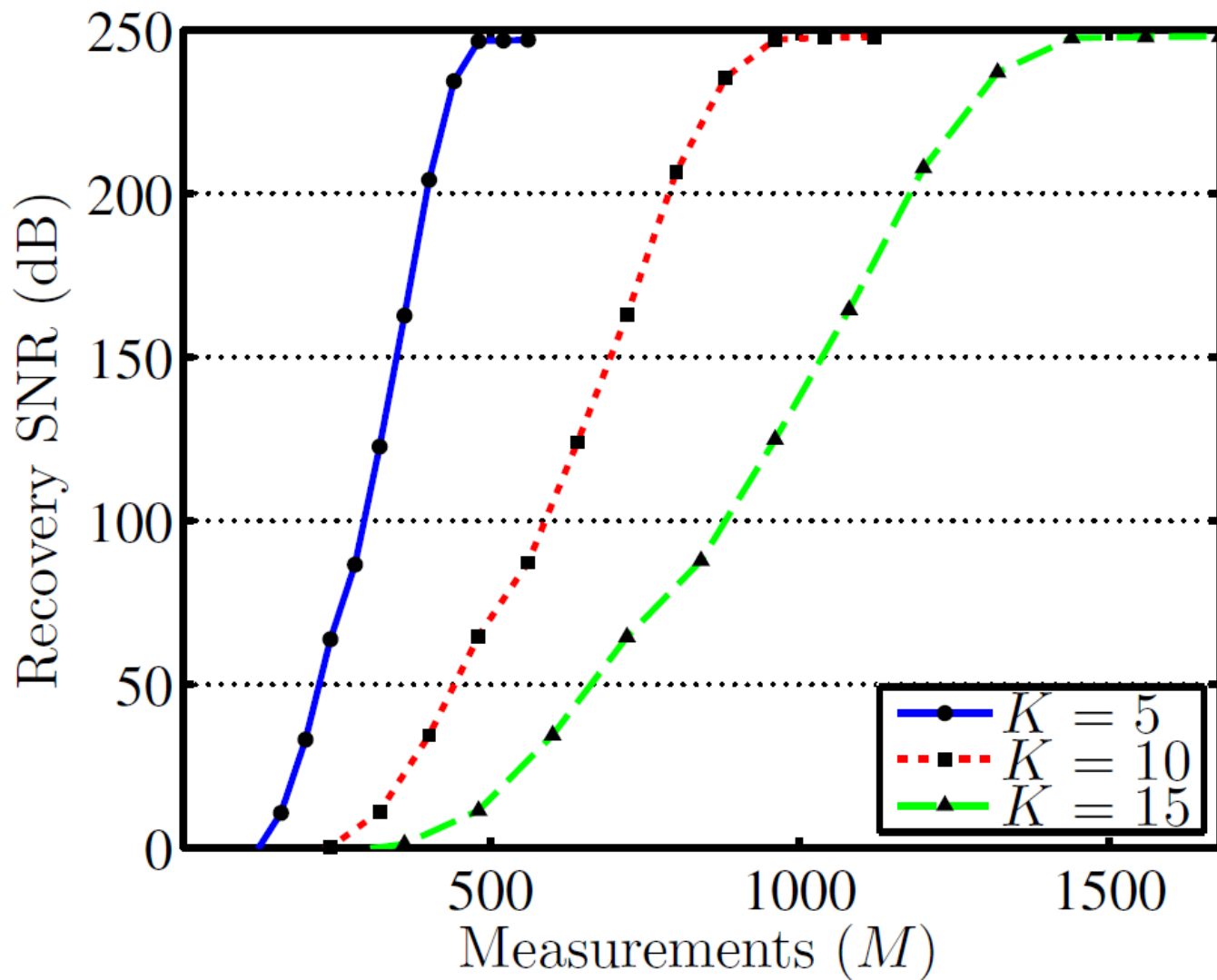
$$N = 4096$$

$$M = 512$$

$$K = 5$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

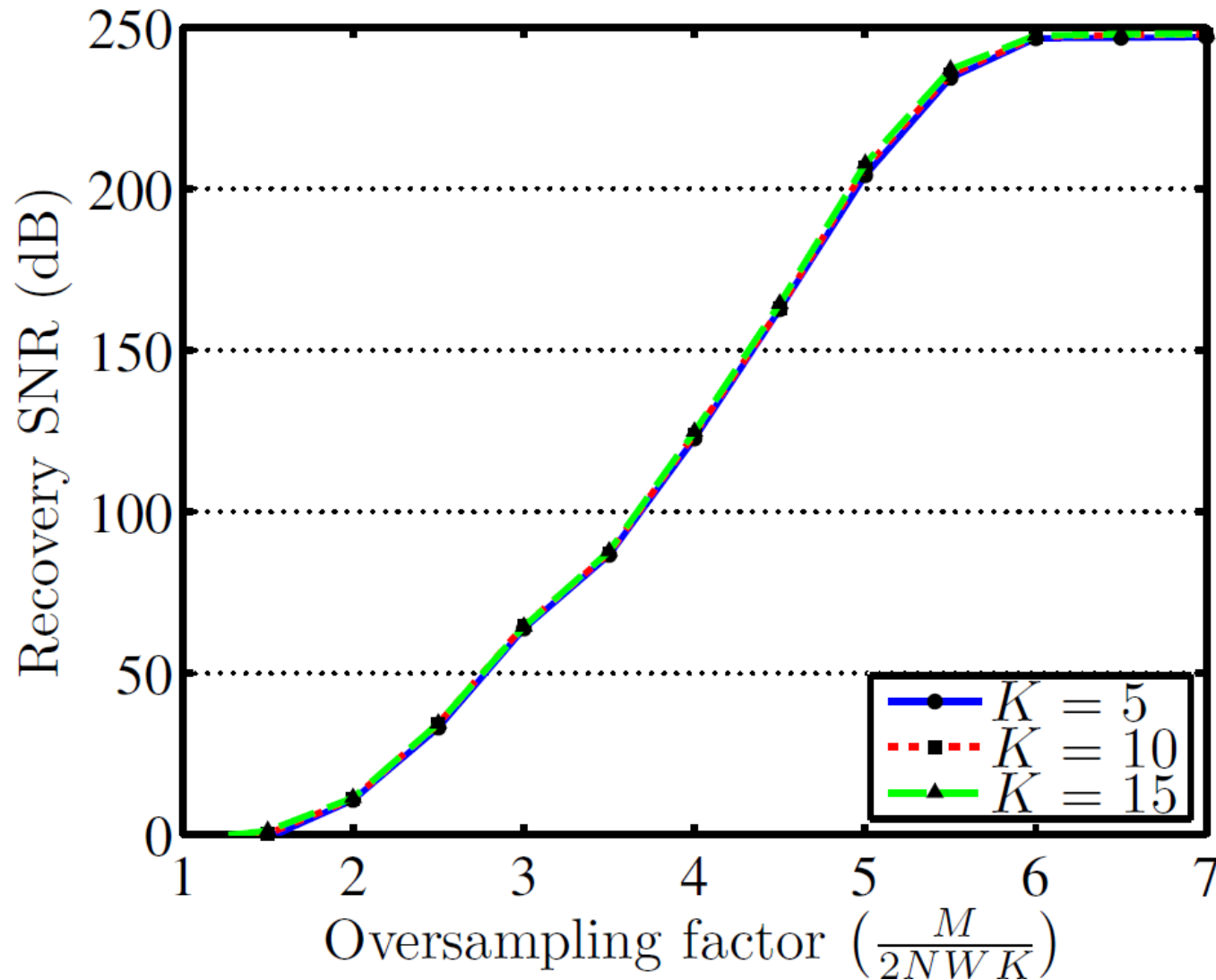
# Empirical Results: Measurements



$$N = 4096$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

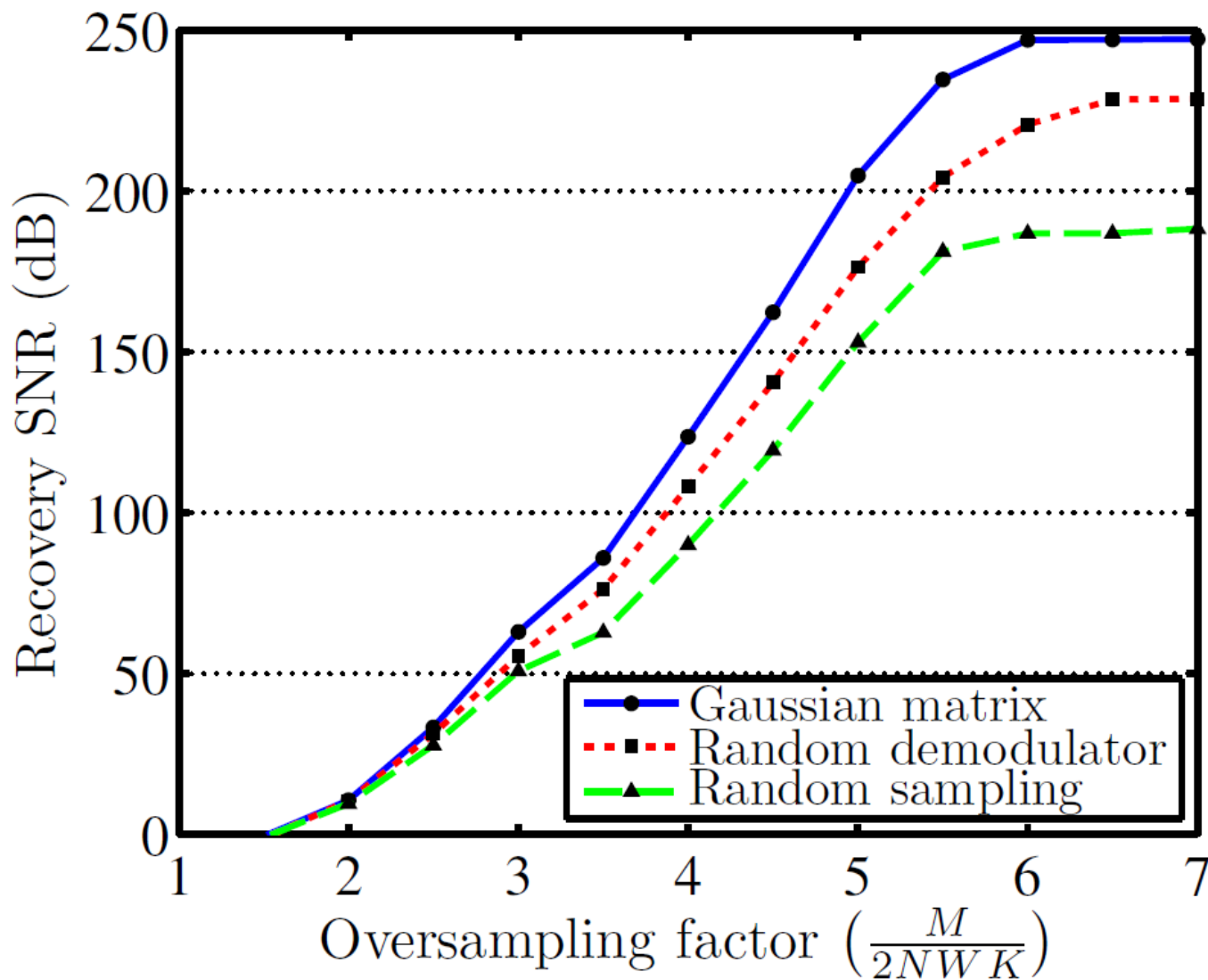
# Empirical Results: Measurements



$$N = 4096$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

# Empirical Results: Real-World Sensors

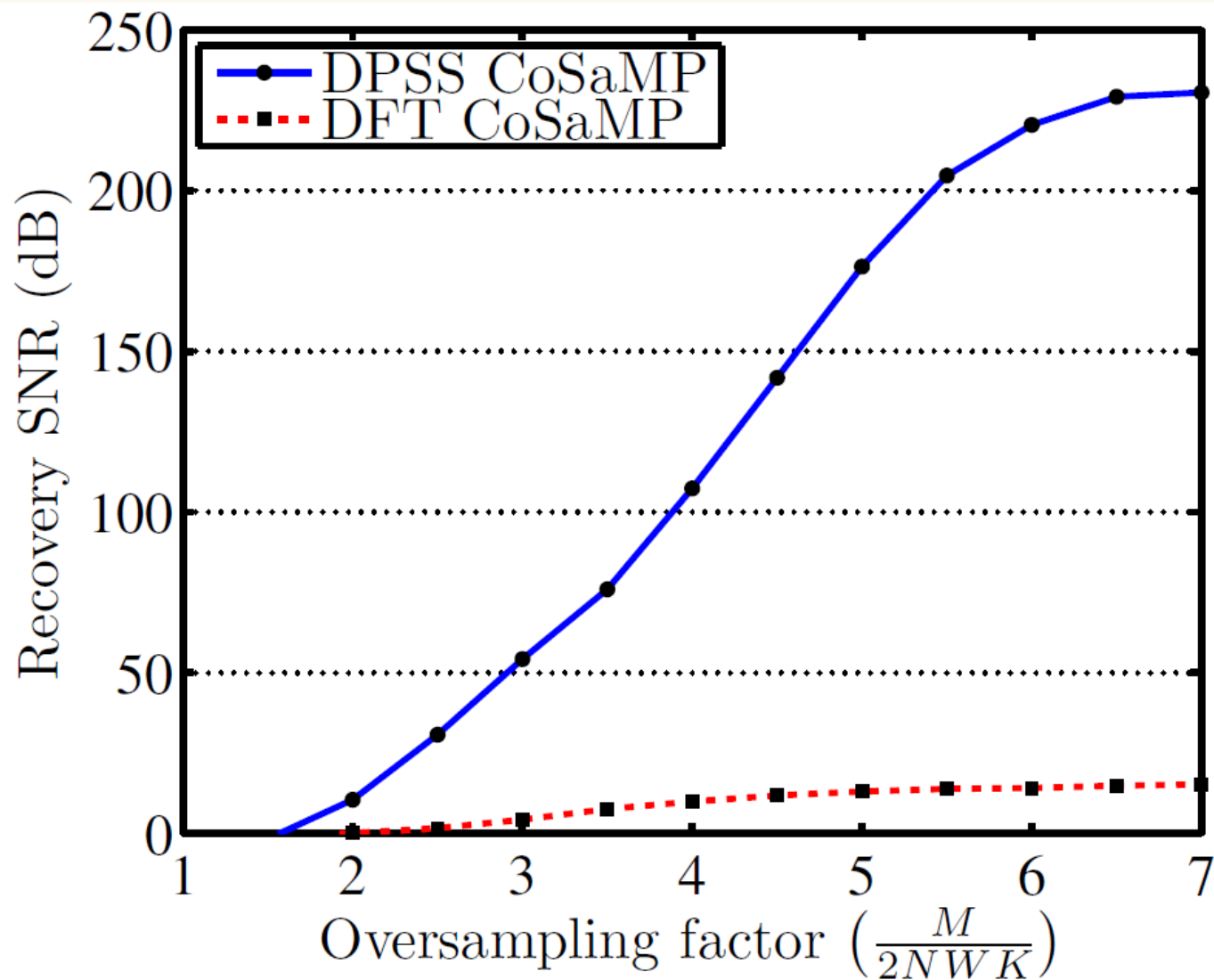


$$N = 4096$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

$$K = 5$$

# Empirical Results: DFT Comparison

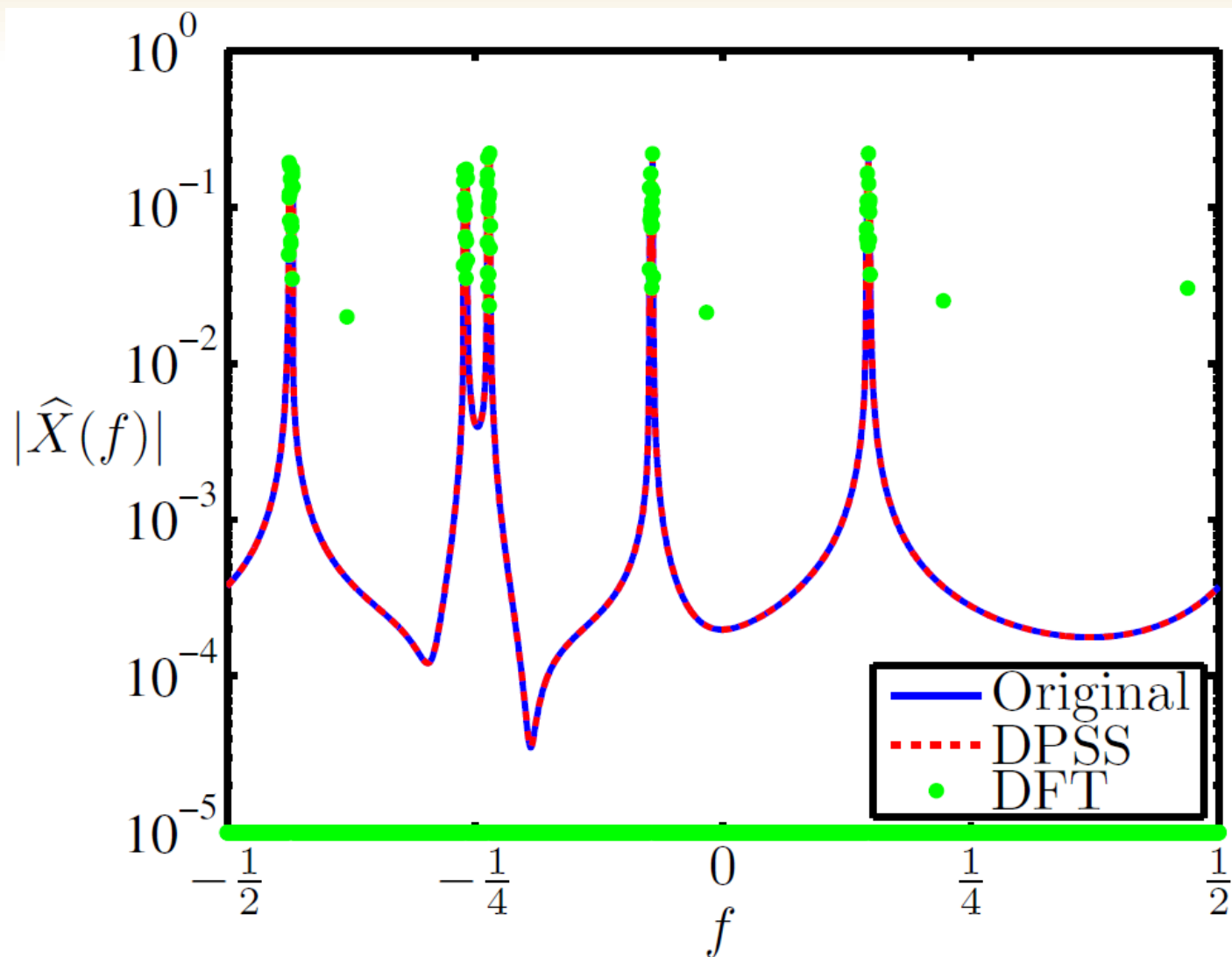


$$N = 4096$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

$$K = 5$$

# Empirical Results: DFT Comparison



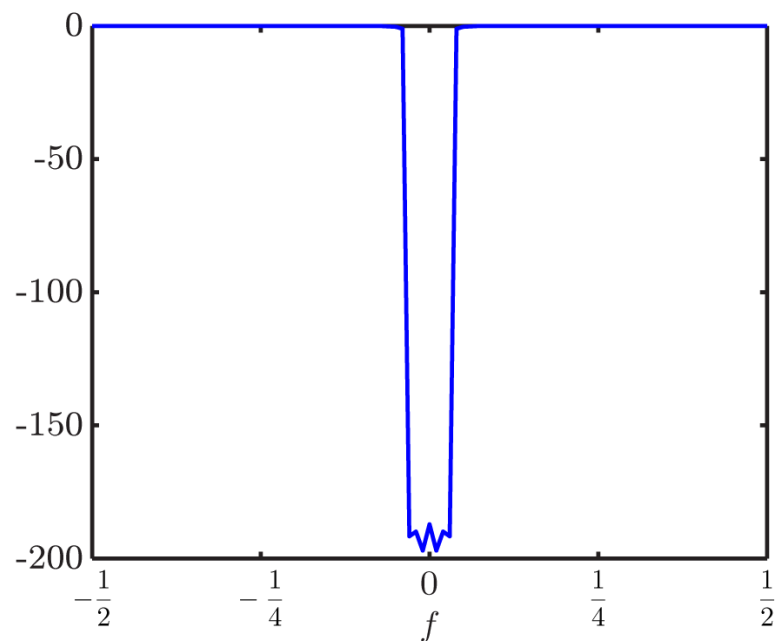
$$N = 4096$$

$$\frac{B}{B_{\text{nyq}}} = \frac{1}{256}$$

$$K = 5$$

# Interference Cancellation

DPSS's can be used to cancel bandlimited interferers *without reconstruction*.



$$P = I - \Phi\Psi_i(\Phi\Psi_i)^\dagger$$

Extremely useful in *compressive signal processing* applications.



# Conclusions

- DPSS's can be used to efficiently represent *most* sampled multiband signals
  - far superior to DFT
- Two types of error: *approximation* + *reconstruction*
  - approximation: small for most signals
  - reconstruction: zero for DPSS-sparse vectors
  - delicate balance in practice, but there is a sweet spot
- This approach combines careful design of  $\Psi$  with more sophisticated sparse models
  - relevant in many contexts beyond ADCs